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The AES-CMAC-96 Algorithm and Its Use with IPsec

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Abstract

The National Institute of Standards and Technology (NIST) has recently specified the Cipher-based Message Authentication Code (CMAC), which is equivalent to the One-Key CBC-MAC1 (OMAC1) algorithm submitted by Iwata and Kurosawa. OMAC1 efficiently reduces the key size of Extended Cipher Block Chaining mode (XCBC). This memo specifies the use of CMAC mode on the authentication mechanism of the IPsec Encapsulating Security Payload (ESP) and the Authentication Header (AH) protocols. This new algorithm is named AES-CMAC-96.

1. Introduction

The National Institute of Standards and Technology (NIST) has recently specified the Cipher-based Message Authentication Code (CMAC). CMAC [NIST-CMAC] is a message authentication code that is based on a symmetric key block cipher such as the Advanced Encryption Standard [NIST-AES]. CMAC is equivalent to the One-Key CBC MAC1 (OMAC1) submitted by Iwata and Kurosawa [OMAC1a, OMAC1b]. OMAC1 is an improvement of the eXtended Cipher Block Chaining mode (XCBC) submitted by Black and Rogaway [XCBCa, XCBCb], which itself is an improvement of the basic CBC-MAC. XCBC efficiently addresses the security deficiencies of CBC-MAC, and OMAC1 efficiently reduces the key size of XCBC.

This memo specifies the usage of CMAC on the authentication mechanism of the IPsec Encapsulating Security Payload [ESP] and Authentication Header [AH] protocols. This new algorithm is named AES-CMAC-96. For further information on AH and ESP, refer to [AH] and [ROADMAP].

2. Basic Definitions

CBC	Cipher Block Chaining mode of operation for message authentication code.
MAC	Message Authentication Code. A bit string of a fixed length, computed by the MAC generation algorithm, that is used to establish the authority and, hence, the integrity of a message.
CMAC	Cipher-based MAC based on an approved symmetric key block cipher, such as the Advanced Encryption Standard.
Key (K)	128-bit (16-octet) key for AES-128 cipher block. Denoted by K.
Message (M)	Message to be authenticated. Denoted by M.
Length (len)	The length of message M in octets. Denoted by len. The minimum value is 0. The maximum value is not specified in this document.
truncate(T,l)	Truncate T (MAC) in most-significant-bit-first (MSB-first) order to a length of l octets.
T	The output of AES-CMAC.

Truncated T The truncated output of AES-CMAC-128 in MSB-first order.

AES-CMAC CMAC generation function based on AES block cipher with 128-bit key.

AES-CMAC-96 IPsec AH and ESP MAC generation function based on AES-CMAC, which truncates the 96 most significant bits of the 128-bit output.

3. AES-CMAC

The core of AES-CMAC-96 is the AES-CMAC [AES-CMAC]. The underlying algorithms for AES-CMAC are the Advanced Encryption Standard cipher block [NIST-AES] and the recently defined CMAC mode of operation [NIST-CMAC]. AES-CMAC provides stronger assurance of data integrity than a checksum or an error detecting code. The verification of a checksum or an error detecting code detects only accidental modifications of the data, while CMAC is designed to detect intentional, unauthorized modifications of the data, as well as accidental modifications. The output of AES-CMAC can validate the input message. Validating the message provides assurance of the integrity and authenticity over the message from the source. According to [NIST-CMAC], at least 64 bits should be used against guessing attacks. AES-CMAC achieves the similar security goal of HMAC [RFC-HMAC]. Since AES-CMAC is based on a symmetric key block cipher (AES), while HMAC is based on a hash function (such as SHA-1), AES-CMAC is appropriate for information systems in which AES is more readily available than a hash function. Detailed information about AES-CMAC is available in [AES-CMAC] and [NIST-CMAC].

4. AES-CMAC-96

For IPsec message authentication on AH and ESP, AES-CMAC-96 should be used. AES-CMAC-96 is a AES-CMAC with 96-bit truncated output in MSB-first order. The output is a 96-bit MAC that will meet the default authenticator length as specified in [AH]. The result of truncation is taken in MSB-first order. For further information on AES-CMAC, refer to [AES-CMAC] and [NIST-CMAC].

Figure 1 describes AES-CMAC-96 algorithm:

In step 1, AES-CMAC is applied to the message M in length len with key K.

In step 2, the output block T is truncated to 12 octets in MSB-first order, and Truncated T (TT) is returned.

```
+++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++
+                    Algorithm AES-CMAC-96                          +
+++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++
+                                                                   +
+   Input    : K (128-bit Key described in Section 4.1)             +
+            : M    (message to be authenticated)                   +
+            : len  (length of message in octets)                   +
+   Output   : Truncated T  (truncated output to length 12 octets)  +
+                                                                   +
+-------------------------------------------------------------------+
+                                                                   +
+   Step 1.  T  := AES-CMAC (K,M,len);                              +
+   Step 2.  TT := truncate (T, 12);                                +
+            return TT;                                             +
+++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++++
```

Figure 1: Algorithm AES-CMAC-96

5. Test Vectors

These test cases are the same as those defined in [NIST-CMAC], with the exception of 96-bit truncation.

```
--------------------------------------------------
K              2b7e1516 28aed2a6 abf71588 09cf4f3c
Subkey Generation
AES_128(key,0) 7df76b0c 1ab899b3 3e42f047 b91b546f
K1             fbeed618 35713366 7c85e08f 7236a8de
K2             f7ddac30 6ae266cc f90bc11e e46d513b

Test Case 1: len = 0
M              <empty string>
AES_CMAC_96    bb1d6929 e9593728 7fa37d12

Test Case 2: len = 16
M              6bc1bee2 2e409f96 e93d7e11 7393172a
AES_CMAC_96    070a16b4 6b4d4144 f79bdd9d

Test Case 3: len = 40
M              6bc1bee2 2e409f96 e93d7e11 7393172a
               ae2d8a57 1e03ac9c 9eb76fac 45af8e51
               30c81c46 a35ce411
AES_CMAC_96    dfa66747 de9ae630 30ca3261

Test Case 4: len = 64
M              6bc1bee2 2e409f96 e93d7e11 7393172a
               ae2d8a57 1e03ac9c 9eb76fac 45af8e51
               30c81c46 a35ce411 e5fbc119 1a0a52ef
               f69f2445 df4f9b17 ad2b417b e66c3710
AES_CMAC_96    51f0bebf 7e3b9d92 fc497417
--------------------------------------------------
```

6. Interaction with the ESP Cipher Mechanism

As of this writing, there are no known issues that preclude the use of AES-CMAC-96 with any specific cipher algorithm.

7. Security Considerations

See the security considerations section of [AES-CMAC].

8. IANA Considerations

The IANA has allocated value 8 for IKEv2 Transform Type 3 (Integrity Algorithm) to the AUTH_AES_CMAC_96 algorithm.

9. Acknowledgements

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10. References

10.1. Normative References

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10.2. Informative References

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