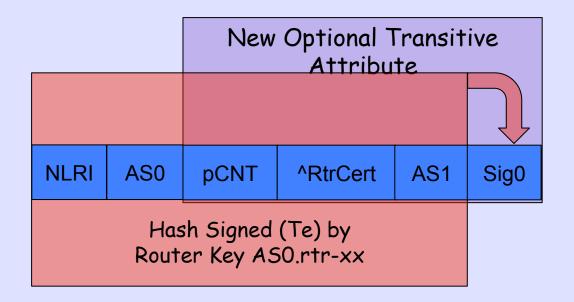
## Estimating CPU Cost of BGPsec on a Router

IEPG / Taipei

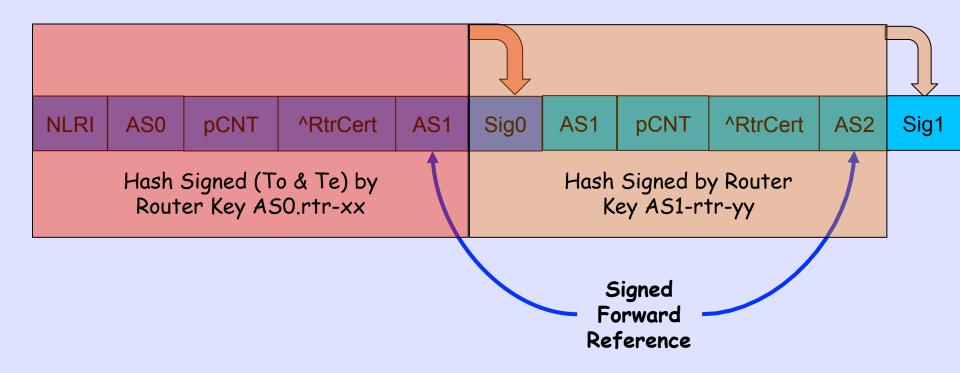
2011.11.13

Kotikalapudi Sriram <kotikalapudi.sriram@nist.gov>
Randy Bush <randy@psg.com>

#### BGPsec from ASO to AS1



#### BGPsec AS1 to AS2

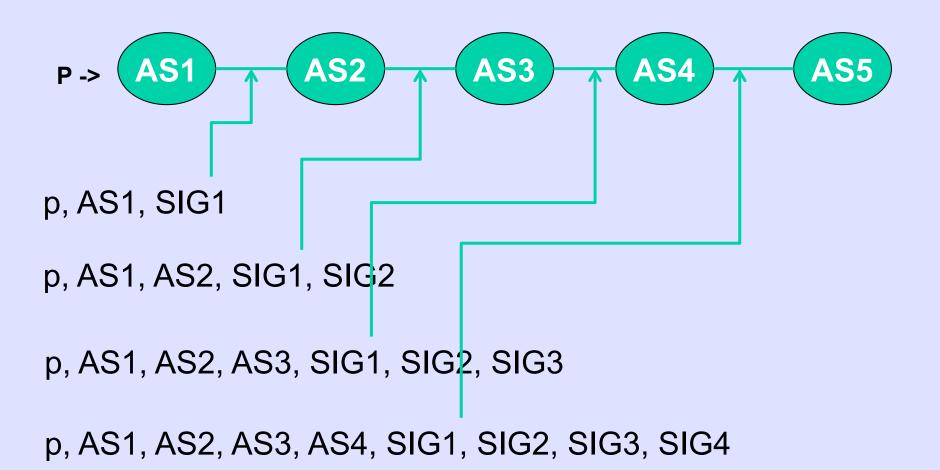


R1 signing over R0's signature is same as signing over entire R0 announcement

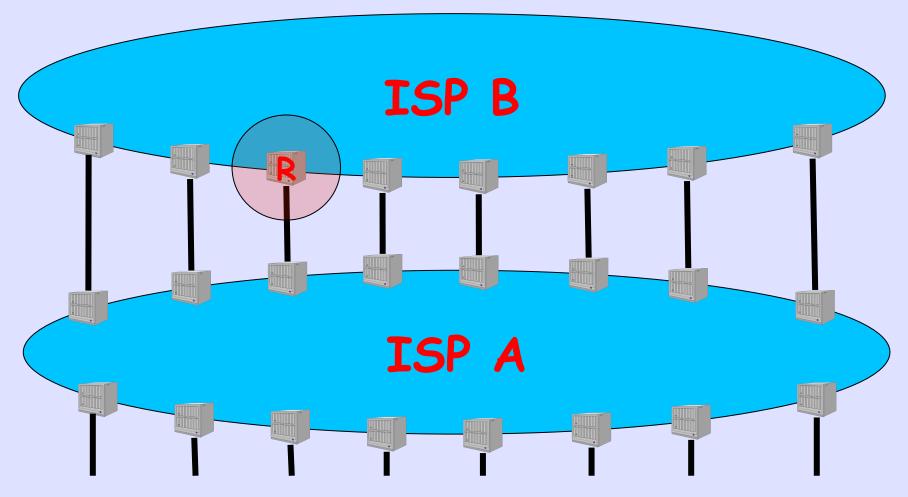
#### BGPsec Islands

- RPKI-Based Origin Validation can be deployed by randomly scattered ISPs
- · Each gets the benefit of origin validation
- · BGPsec depends on your neighbor signing
- It will deploy as islands which eventually interconnect

#### We Draw Pictures Like This



## But Reality is This



A's Customer 'Cone'

#### Number of Paths

- One ISP router, R, has many paths for prefix P
- All but one are from iBGP peers
- BGPsec spec says R does not validate paths received from iBGP peers
- I.e. R has to validate only one path for each P from peer A

#### Some Largish ISPs Cones

#### Very Large Global

- 1 1353 --- **ISP's Own Pf**x
- 2 21586 --- **BGP Cust Pfx**
- 3 6820 --- Cust's Cust Pfx
- 4 1627 --- ...
- 5 942
- 6 45
- 7 14
- 8 6

#### Very Large Global

- 1 620
- 2 16028
- 3 9434
- 4 2922
- 5 435
- 6 46
- 7 15
- 8 27
- 9 1

#### Large Global

- 1 443
- 2 8197
- 3 8052
- 4 2715
- 5 387
- 6 37
- 7 48
- 8 157
- 9 2

#### Large Global

- 1 501
- 2 3686
- 3 3603
- 4 816
- 5 45
- 6 9
- 8 1

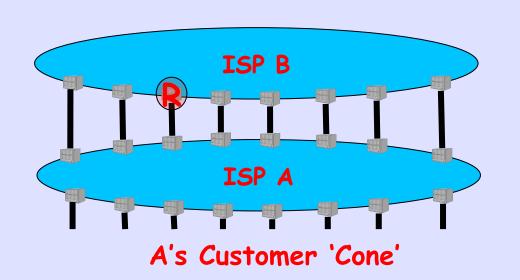
#### Asian Regional

- 1 152
- 2 791
- 3 120
- 4 35
- 5 3 # pfxs path length

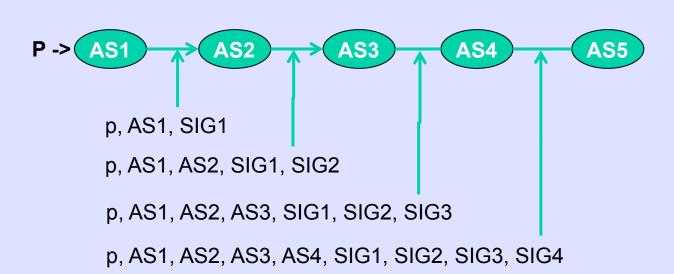
#### Yes, there are rather long tails

Yes, we removed prepending

#### Incremental Deployment



If A and B Deploy BGPsec, What is the Load on a Router?



Now this Picture Makes Sense!

# Cost to Sign/Validate Using One Core

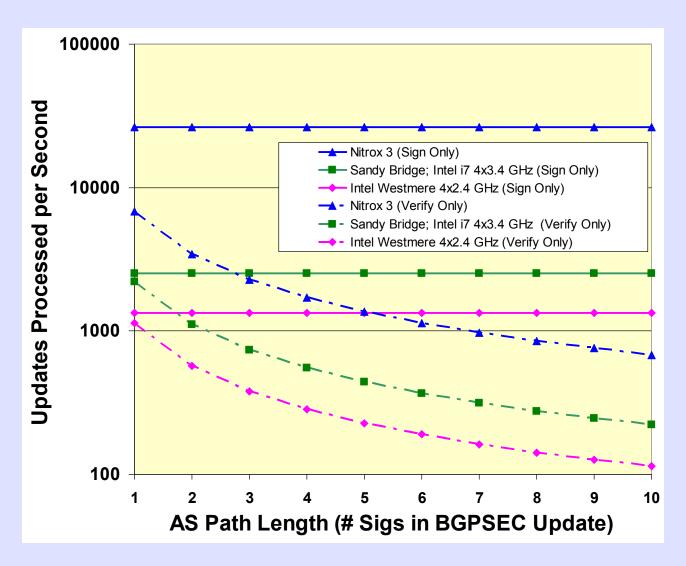
	Operations per second					
			amd64, Sandy			
		/	Bridge; 2011			
		amd64; Westmere	Intel i7-	NITROX PX PCI-	NITROX III PCI-	
	Intel Core 2 Duo,	(206c2); 2010 Intel	2600K;	Express CN1620 -	Express CNN3570-	
	64-bit, 3 GHz,	Xeon E5620; 4 x	3400MHz;	P¢le Look-aside	PCle Look-aside	
	8GB, Linux 5.7	2400MHz	threads	Processor	Processor	
ECDSA-P256 Verify	890	1139	2215	854	6832	
ECDSA-P256 Sign	1100	1335	2530	3293	26344	

 Source: eBACS: ECRYPT Benchmarking of Cryptographic Systems

http://bench.cr.yp.to/results-sign.htm

And: Cavium, Inc. (private communication)

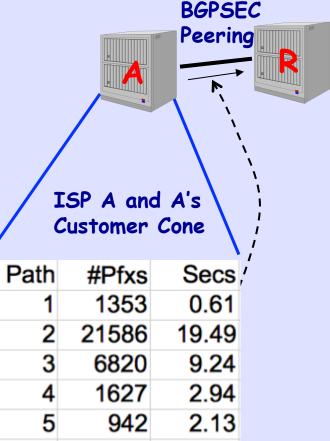
#### Updates Per Second



#### Validation Cost Model

**BGPSEC** 

Peering



45

14

0.12

0.04

0.02

34.59

CPU Time on R if Session to C is Reset

8 6 Total Seconds

6

CPU Time

Session to

A is Reset

on R if

Path	#Ptxs	Secs
1	620	0.28
2	16028	14.47
3	9434	12.78
4	2922	5.28
5	435	0.98
6	46	0.12
7	15	0.05
8	27	0.10
9	1	0.00
Total Se	34.06	

ISP C and C's

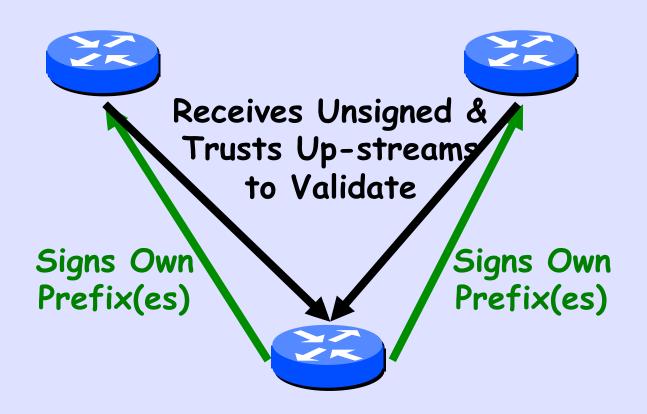
Customer Cone

**4Df.** 

## Signing Cost

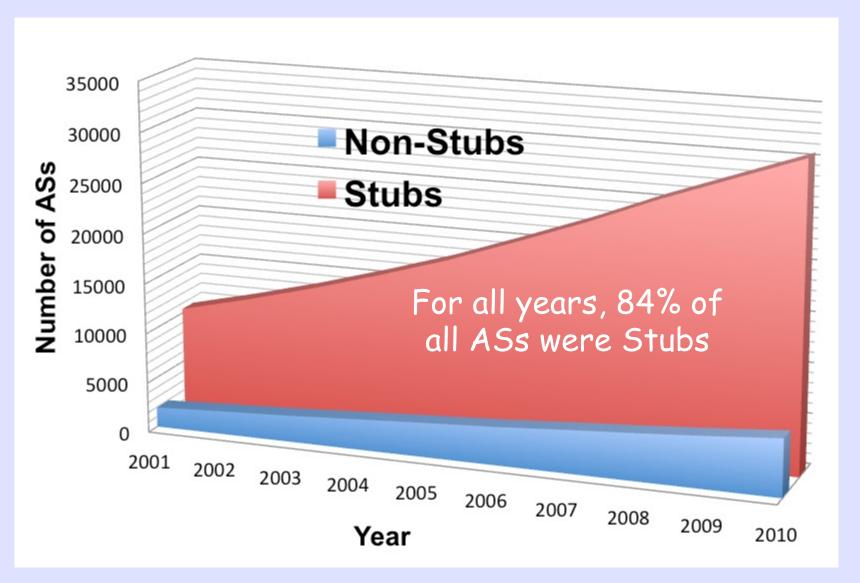
- You only sign once, irrespective of path length
- You only sign toward BGPsec speakers
- Though the cost of stripping
   BGPsec toward non-speakers may
   be on the order of signing

## Need not Sign To Stubs



Only Needs to Have Own Private Key, No Other Crypto or RPKI Data No Hardware Upgrade!!

#### Stub ASs vs Transit



## BGP Peers per Router

ISP	BGP Peers	BGP Custs
W	29	95
X	3-4	20
У	6	12
Z	8	16

These numbers are from real ISPs, but large ones

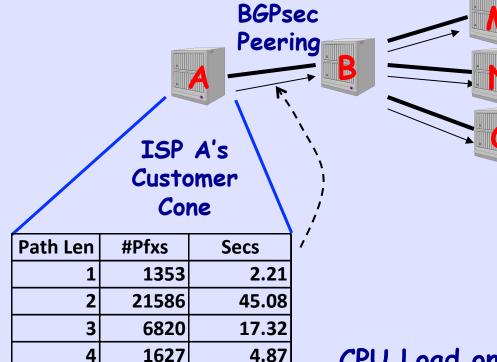
### Signing Bottom Line

- Except for W, it comes to 2-3 BGPsec customers per aggregation router
- Say 400k routes at 2530 sigs/sec
- (3\*400000)/2530 = 475 seconds
- But this presumes the entire Internet is signed, which is a looooooooong time from now
- · But W will eventually have a problem!

2011.11.13 BGPsec Crypto Cost

17

#### CPU for Validation and Signing



3.24

0.18

0.06

0.03

**73** 

942

45

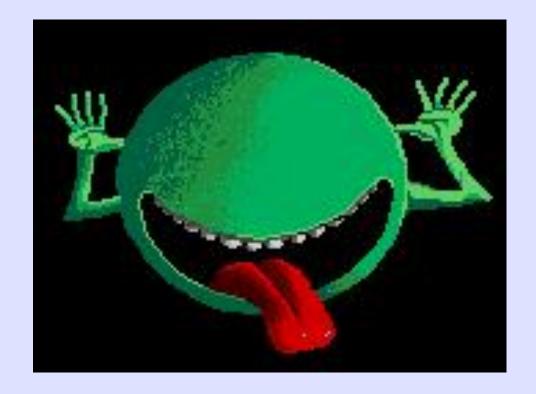
14

Total (seconds)

6

- · B peers with four BGPsec peers
- B's other peers are not BGPsec aware

CPU Load on B, including Validation & Signing, if Session to A is Reset.



## So Don't Panic, Engineer Prudently