Network Working Group Request for Comments: 4898 Category: Standards Track

TCP Extended Statistics MIB

Status of This Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

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Abstract

This document describes extended performance statistics for TCP. They are designed to use TCP's ideal vantage point to diagnose performance problems in both the network and the application. If a network-based application is performing poorly, TCP can determine if the bottleneck is in the sender, the receiver, or the network itself. If the bottleneck is in the network, TCP can provide specific information about its nature.

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1. Introduction

This document describes extended performance statistics for TCP. They are designed to use TCP's ideal vantage point to diagnose performance problems in both the network and the application. If a network-based application is performing poorly, TCP can determine if the bottleneck is in the sender, the receiver, or the network itself. If the bottleneck is in the network, TCP can provide specific information about its nature.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119.

The Simple Network Management Protocol (SNMP) objects defined in this document extend TCP MIB, as specified in RFC 4022 [RFC4022]. In addition to several new scalars and other objects, it augments two tables and makes one clarification to RFC 4022. Existing management stations for the TCP MIB are expected to be fully compatible with these clarifications.

2. The Internet-Standard Management Framework

For a detailed overview of the documents that describe the current Internet-Standard Management Framework, please refer to section 7 of RFC 3410 [RFC3410].

Managed objects are accessed via a virtual information store, termed the Management Information Base or MIB. MIB objects are generally accessed through the Simple Network Management Protocol (SNMP). Objects in the MIB are defined using the mechanisms defined in the Structure of Management Information (SMI). This memo specifies a MIB module that is compliant to the SMIv2, which is described in STD 58, RFC 2578 [RFC2578], STD 58, RFC 2579 [RFC2579] and STD 58, RFC 2580 [RFC2580].

3. Overview

The TCP-ESTATS-MIB defined in this memo consists of two groups of scalars, seven tables, and two notifications:

* The first group of scalars contain statistics of the TCP protocol engine not covered in RFC 4022. This group consists of the single scalar tcpEStatsListenerTableLastChange, which provides management stations with an easier mechanism to validate their listener caches.

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- * The second group of scalars consist of knobs to enable and disable information collection by the tables containing connection-related statistics/information. For example, the tcpEStatsControlPath object controls the activation of the tcpEStatsPathTable. The tcpEStatsConnTableLatency object determines how long connection table rows are retained after a TCP connection transitions into the closed state.
- * The tcpEStatsListenerTable augments tcpListenerTable in TCP-MIB [RFC4022] to provide additional information on the active TCP listeners on a device. It supports objects to monitor and diagnose SYN-flood denial-of-service attacks as described below.
- * The tcpEStatsConnectIdTable augments the tcpConnectionTable in TCP-MIB [RFC4022] to provide a mapping between connection 4-tuples (which index tcpConnectionTable) and an integer connection index, tcpEStatsConnectIndex. The connection index is used to index into the five remaining tables in this MIB module, and is designed to facilitate rapid polling of multiple objects associated with one TCP connection.
- * The tcpEStatsPerfTable contains objects that are useful for measuring TCP performance and first check problem diagnosis.
- * The tcpEStatsPathTable contains objects that can be used to infer detailed behavior of the Internet path, such as the extent that there are segment losses or reordering, etc.
- * The tcpEStatsStackTable contains objects that are most useful for determining how well the TCP control algorithms are coping with this particular path.
- * The tcpEStatsAppTable provides objects that are useful for determining if the application using TCP is limiting TCP performance.
- * The tcpEStatsTuneTable provides per-connection controls that can be used to work around a number of common problems that plague TCP over some paths.
- * The two notifications defined in this MIB module are tcpEStatsEstablishNotification, indicating that a new connection has been accepted (or established, see below), and tcpEStatsCloseNotification, indicating that an existing connection has recently closed.

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3.1. MIB Initialization and Persistence

The TCP protocol itself is specifically designed not to preserve any state whatsoever across system reboots, and enforces this by requiring randomized Initial Sequence numbers and ephemeral ports under any conditions where segments from old connections might corrupt new connections following a reboot.

All of the objects in the MIB MUST have the same persistence properties as the underlying TCP implementation. On a reboot, all zero-based counters MUST be cleared, all dynamically created table rows MUST be deleted, and all read-write objects MUST be restored to their default values. It is assumed that all TCP implementation have some initialization code (if nothing else, to set IP addresses) that has the opportunity to adjust tcpEStatsConnTableLatency and other read-write scalars controlling the creation of the various tables, before establishing the first TCP connection. Implementations MAY also choose to make these control scalars persist across reboots.

The ZeroBasedCounter32 and ZeroBasedCounter64 objects in the listener and connection tables are initialized to zero when the table row is created.

The tcpEStatsConnTableLatency object determines how long connection table rows are retained after a TCP connection transitions into the closed state, to permit reading final connection completion statistics. In RFC 4022 (TCP-MIB), the discussion of tcpConnectionTable row latency (page 9) the words "soon after" are understood to mean after tcpEStatsConnTableLatency, such that all rows of all tables associated with one connection are retained at least tcpEStatsConnTableLatency after connection close. This clarification to RFC 4022 only applies when TCP-ESTATS-MIB is implemented. If TCP-ESTATS-MIB is not implemented, RFC 4022 permits an unspecified delay between connection close and row deletion.

3.2. Relationship to TCP Standards

There are more than 70 RFCs and other documents that specify various aspects of the Transmission Control Protocol (TCP) [RFC4614]. While most protocols are completely specified in one or two documents, this has not proven to be feasible for TCP. TCP implements a reliable end-to-end data transport service over a very weakly constrained IP datagram service. The essential problem that TCP has to solve is balancing the applications need for fast and reliable data transport against the need to make fair, efficient, and equitable use of network resources, with only sparse information about the state of the network or its capabilities.

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TCP maintains this balance through the use of many estimators and heuristics that regulate various aspects of the protocol. For example, RFC 2988 describes how to calculate the retransmission timer (RTO) from the average and variance of the network round-trip-time (RTT), as estimated from the round-trip time sampled on some data segments. Although these algorithms are standardized, they are a compromise which is optimal for only common Internet environments. Other estimators might yield better results (higher performance or more efficient use of the network) in some environments, particularly under uncommon conditions.

It is the consensus of the community that nearly all of the estimators and heuristics used in TCP might be improved through further research and development. For this reason, nearly all TCP documents leave some latitude for future improvements, for example, by the use of "SHOULD" instead of "MUST" [RFC2119]. Even standard algorithms that are required because they critically effect fairness or the dynamic stability of Internet congestion control, include some latitude for evolution. As a consequence, there is considerable diversity in the details of the TCP implementations actually in use today.

The fact that the underlying algorithms are not uniform makes it difficult to tightly specify a MIB. We could have chosen the point of view that the MIB should publish precisely defined metrics of the network path, even if they are different from the estimators in use by TCP. This would make the MIB more useful as a measurement tool, but less useful for understanding how any specific TCP implementation is interacting with the network path and upper protocol layers. We chose instead to have the MIB expose the estimators and important states variables of the algorithms in use, without constraining the TCP implementation.

As a consequence, the MIB objects are defined in terms of fairly abstract descriptions (e.g., round-trip time), but are intended to expose the actual estimators or other state variables as they are used in TCP implementations, possibly transformed (e.g., scaled or otherwise adjusted) to match the spirit of the object descriptions in this document.

This may mean that MIB objects may not be exactly comparable between two different TCP implementations. A general management station can only assume the abstract descriptions, which are useful for a general assessment of how TCP is functioning. To a TCP implementer with detailed knowledge about the TCP implementation on a specific host, this MIB might be useful for debugging or evaluating the algorithms in their implementation.

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Under no conditions is this MIB intended to constrain TCP to use (or exclude) any particular estimator, heuristic, algorithm, or implementation.

3.3. Diagnosing SYN-Flood Denial-of-Service Attacks

The tcpEStatsListenerTable is specifically designed to provide information that is useful for diagnosing SYN-flood Denial-of-Service attacks, where a server is overwhelmed by forged or otherwise malicious connection attempts. There are several different techniques that can be used to defend against SYN-flooding but none are standardized [Edd06]. These different techniques all have the same basic characteristics that are instrumentable with a common set of objects, even though the techniques differ greatly in the details.

All SYN-flood defenses avoid allocating significant resources (memory or CPU) to incoming (passive open) connections until the connections meet some liveness criteria (to defend against forged IP source addresses) and the server has sufficient resources to process the incoming request. Note that allocating resources is an implementation-specific event that may not correspond to an observable protocol event (e.g., segments on the wire). There are two general concepts that can be applied to all known SYN-flood defenses. There is generally a well-defined event when a connection is allocated full resources, and a "backlog" -- a queue of embryonic connections that have been allocated only partial resources.

In many implementations, incoming TCP connections are allocated resources as a side effect of the POSIX [POSIX] accept() call. For this reason we use the terminology "accepting a connection" to refer to this event: committing sufficient network resources to process the incoming request. Accepting a connection typically entails allocating memory for the protocol control block [RFC793], the per-connection table rows described in this MIB and CPU resources, such as process table entries or threads.

Note that it is not useful to accept connections before they are ESTABLISHED, because this would create an easy opportunity for Denial-of-Service attacks, using forged source IP addresses.

The backlog consists of connections that are in SYN-RCVD or ESTABLISHED states, that have not been accepted. For purposes of this MIB, we assume that these connections have been allocated some resources (e.g., an embryonic protocol control block), but not full resources (e.g., do not yet have MIB table rows).

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Note that some SYN-Flood defenses dispense with explicit SYN-RCVD state by cryptographically encoding the state in the ISS (initial sequence number sent) of the SYN-ACK (sometimes called a syn-cookie), and then using the sequence number of the first ACK to reconstruct the SYN-RCVD state before transitioning to the ESTABLISHED state. For these implementations there is no explicit representation of the SYN-RCVD state, and the backlog only consists of connections that are ESTABLISHED and are waiting to be ACCEPTED.

Furthermore, most SYN-flood defenses have some mechanism to throttle connections that might otherwise overwhelm this endpoint. They generally use some combination of discarding incoming SYNs and discarding connections already in the backlog. This does not cause all connections from legitimate clients to fail, as long as the clients retransmit the SYN or first ACK as specified in RFC 793. Most diversity in SYN flood defenses arise from variations in these algorithms to limit load, and therefore cannot be instrumented with a common standard MIB.

The Listen Table instruments all passively opened TCP connections in terms of observable protocol events (e.g., sent and received segments) and resource allocation events (entering the backlog and being accepted). This approach eases generalization to SYN-flood mechanisms that use alternate TCP state transition diagrams and implicit mechanisms to encode some states.

4. TCP Extended Statistics MIB

This MIB module IMPORTS definitions from [RFC2578], [RFC2579], [RFC2580], [RFC2856], [RFC4022], and [RFC4502]. It uses REFERENCE clauses to refer to [RFC791], [RFC793], [RFC1122], [RFC1191], [RFC1323], [RFC2018], [RFC2581], [RFC2861], [RFC2883], [RFC2988], [RFC3168], [RFC3260], [RFC3517], [RFC3522], and [RFC3742].

TCP-ESTATS-MIB DEFINITIONS ::= BEGIN IMPORTS MODULE-IDENTITY, Counter32, Integer32, Unsigned32, Gauge32, OBJECT-TYPE, mib-2, NOTIFICATION-TYPE FROM SNMPv2-SMI -- [RFC2578] MODULE-COMPLIANCE, OBJECT-GROUP, NOTIFICATION-GROUP FROM SNMPv2-CONF -- [RFC2580] ZeroBasedCounter32 FROM RMON2-MIB -- [RFC4502] ZeroBasedCounter64 FROM HCNUM-TC -- [RFC2856] TEXTUAL-CONVENTION, DateAndTime, TruthValue, TimeStamp

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FROM SNMPv2-TC -- [RFC2579] tcpListenerEntry, tcpConnectionEntry FROM TCP-MIB; -- [RFC4022] tcpEStatsMIB MODULE-IDENTITY LAST-UPDATED "200705180000Z" -- 18 May 2007 ORGANIZATION "IETF TSV Working Group" CONTACT-INFO "Matt Mathis John Heffner Web100 Project Pittsburgh Supercomputing Center 300 S. Craig St. Pittsburgh, PA 15213 Email: mathis@psc.edu, jheffner@psc.edu Rajiv Raghunarayan Cisco Systems Inc. San Jose, CA 95134 Phone: 408 853 9612 Email: raraghun@cisco.com Jon Saperia 84 Kettell Plain Road Stow, MA 01775 Phone: 617-201-2655 Email: saperia@jdscons.com " DESCRIPTION "Documentation of TCP Extended Performance Instrumentation variables from the Web100 project. [Web100] All of the objects in this MIB MUST have the same persistence properties as the underlying TCP implementation. On a reboot, all zero-based counters MUST be cleared, all

dynamically created table rows MUST be deleted, and all read-write objects MUST be restored to their default values.

It is assumed that all TCP implementation have some initialization code (if nothing else to set IP addresses) that has the opportunity to adjust tcpEStatsConnTableLatency and other read-write scalars controlling the creation of the various tables, before establishing the first TCP connection. Implementations MAY also choose to make these control scalars persist across reboots.

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```
REVISION "200705180000Z" -- 18 May 2007
DESCRIPTION
"Initial version, published as RFC 4898."
::= { mib-2 156 }
tcpEStatsNotifications OBJECT IDENTIFIER ::= { tc
tcpEStatsMIBObjects OBJECT IDENTIFIER ::= { tc
```

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```
tcpEStatsNotifications OBJECT IDENTIFIER ::= { tcpEStatsMIB 0 }
tcpEStatsMIBObjects OBJECT IDENTIFIER ::= { tcpEStatsMIB 1 }
tcpEStatsConformance OBJECT IDENTIFIER ::= { tcpEStatsMIB 2 }
tcpEStatsOBJECT IDENTIFIER ::= { tcpEStatsMIBObjects 1 }tcpEStatsControlOBJECT IDENTIFIER ::= { tcpEStatsMIBObjects 2 }tcpEStatsScalarOBJECT IDENTIFIER ::= { tcpEStatsMIBObjects 3 }
_ _
-- Textual Conventions
_ _
TcpEStatsNegotiated ::= TEXTUAL-CONVENTION
   STATUS
               current
   DESCRIPTION
        "Indicates if some optional TCP feature was negotiated.
         Enabled(1) indicates that the feature was successfully
         negotiated on, which generally requires both hosts to agree
         to use the feature.
         selfDisabled(2) indicates that the local host refused the
         feature because it is not implemented, configured off, or
         refused for some other reason, such as the lack of
         resources.
         peerDisabled(3) indicates that the local host was willing
         to negotiate the feature, but the remote host did not
         do so."
   SYNTAX INTEGER {
                  enabled(1),
                 selfDisabled(2),
                 peerDisabled(3)
         }
-- TCP Extended statistics scalars
tcpEStatsListenerTableLastChange OBJECT-TYPE
    SYNTAX TimeStamp
    MAX-ACCESS read-only
    STATUS current
    DESCRIPTION
```

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"The value of sysUpTime at the time of the last creation or deletion of an entry in the tcpListenerTable. If the number of entries has been unchanged since the last re-initialization of the local network management subsystem, then this object contains a zero value." ::= { tcpEStatsScalar 3 } _ _ -- The tcpEStatsControl Group _ _ -- The scalar objects in this group are used to control the -- activation and deactivation of the TCP Extended Statistics -- tables and notifications in this module. _ _ tcpEStatsControlPath OBJECT-TYPE SYNTAXTruthValueMAX-ACCESSread-writeCTATUScumment current STATUS DESCRIPTION "Controls the activation of the TCP Path Statistics table. A value 'true' indicates that the TCP Path Statistics table is active, while 'false' indicates that the table is inactive." DEFVAL { false } ::= { tcpEStatsControl 1 } tcpEStatsControlStack OBJECT-TYPE SYNTAX TruthValue MAX-ACCESS read-write STATUS current DESCRIPTION "Controls the activation of the TCP Stack Statistics table. A value 'true' indicates that the TCP Stack Statistics table is active, while 'false' indicates that the table is inactive." DEFVAL { false } ::= { tcpEStatsControl 2 } tcpEStatsControlApp OBJECT-TYPE SYNTAX TruthValue MAX-ACCESS read-write

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STATUS current DESCRIPTION "Controls the activation of the TCP Application Statistics table. A value 'true' indicates that the TCP Application Statistics table is active, while 'false' indicates that the table is inactive." DEFVAL { false } ::= { tcpEStatsControl 3 } tcpEStatsControlTune OBJECT-TYPE SYNTAX TruthValue MAX-ACCESS read-write STATUS current DESCRIPTION "Controls the activation of the TCP Tuning table. A value 'true' indicates that the TCP Tuning table is active, while 'false' indicates that the table is inactive." DEFVAL { false } ::= { tcpEStatsControl 4 } tcpEStatsControlNotify OBJECT-TYPE SYNTAX TruthValue MAX-ACCESS read-write STATUS current DESCRIPTION "Controls the generation of all notifications defined in this MIB. A value 'true' indicates that the notifications are active, while 'false' indicates that the notifications are inactive." DEFVAL { false } ::= { tcpEStatsControl 5 } tcpEStatsConnTableLatency OBJECT-TYPE SYNTAX Unsigned32 UNITS "seconds" MAX-ACCESS read-write STATUS current DESCRIPTION "Specifies the number of seconds that the entity will retain entries in the TCP connection tables, after the connection first enters the closed state. The entity SHOULD provide a configuration option to enable

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```
customization of this value. A value of 0
           results in entries being removed from the tables as soon as
           the connection enters the closed state. The value of
           this object pertains to the following tables:
             tcpEStatsConnectIdTable
              tcpEStatsPerfTable
             tcpEStatsPathTable
             tcpEStatsStackTable
             tcpEStatsAppTable
             tcpEStatsTuneTable"
     DEFVAL \{0\}
     ::= { tcpEStatsControl 6 }
_ _
-- Listener Table
_ _
tcpEStatsListenerTable OBJECT-TYPE
     SYNTAX SEQUENCE OF TCPEStatsListenerEntry
    MAX-ACCESS not-accessible
     STATUS current
     DESCRIPTION
          "This table contains information about TCP Listeners,
          in addition to the information maintained by the
          tcpListenerTable RFC 4022."
     ::= { tcpEStats 1 }
tcpEStatsListenerEntry OBJECT-TYPE
     SYNTAX TcpEStatsListenerEntry
    MAX-ACCESS not-accessible
    STATUS current
    DESCRIPTION
          "Each entry in the table contains information about
          a specific TCP Listener."
     AUGMENTS { tcpListenerEntry }
     ::= { tcpEStatsListenerTable 1 }
TcpEStatsListenerEntry ::= SEQUENCE {
         EstistenerEntry::=SEQUENCE {tcpEStatsListenerStartTimeTimeStamp,tcpEStatsListenerSynRcvdZeroBasedCounter32,tcpEStatsListenerInitialZeroBasedCounter32,tcpEStatsListenerEstablishedZeroBasedCounter32,tcpEStatsListenerAcceptedZeroBasedCounter32,tcpEStatsListenerExceedBacklogZeroBasedCounter32,tcpEStatsListenerHCSynRcvdZeroBasedCounter64,tcpEStatsListenerHCInitialZeroBasedCounter64,tcpEStatsListenerHCInitialZeroBasedCounter64,
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tcpEStatsListenerHCAccepted ZeroBasedCounter64, tcpEStatsListenerHCExceedBacklog ZeroBasedCounter64, tcpEStatsListenerCurConnsGauge32,tcpEStatsListenerMaxBacklogUnsigned32,tcpEStatsListenerCurBacklogGauge32, Gauge32, tcpEStatsListenerCurEstabBacklog Gauge32 } tcpEStatsListenerStartTime OBJECT-TYPE SYNTAX TimeStamp MAX-ACCESS read-only STATUS current DESCRIPTION "The value of sysUpTime at the time this listener was established. If the current state was entered prior to the last re-initialization of the local network management subsystem, then this object contains a zero value." ::= { tcpEStatsListenerEntry 1 } tcpEStatsListenerSynRcvd OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of SYNs which have been received for this listener. The total number of failed connections for all reasons can be estimated to be tcpEStatsListenerSynRcvd minus tcpEStatsListenerAccepted and tcpEStatsListenerCurBacklog." ::= { tcpEStatsListenerEntry 2 } tcpEStatsListenerInitial OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The total number of connections for which the Listener has allocated initial state and placed the connection in the backlog. This may happen in the SYN-RCVD or ESTABLISHED states, depending on the implementation." ::= { tcpEStatsListenerEntry 3 } tcpEStatsListenerEstablished OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION

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```
"The number of connections that have been established to
       this endpoint (e.g., the number of first ACKs that have
       been received for this listener)."
   ::= { tcpEStatsListenerEntry 4 }
tcpEStatsListenerAccepted
                          OBJECT-TYPE
  SYNTAX ZeroBasedCounter32
  MAX-ACCESS read-only
  STATUS current
  DESCRIPTION
     "The total number of connections for which the Listener
      has successfully issued an accept, removing the connection
      from the backlog."
   ::= { tcpEStatsListenerEntry 5 }
tcpEStatsListenerExceedBacklog OBJECT-TYPE
  SYNTAX ZeroBasedCounter32
  MAX-ACCESS read-only
  STATUS current
  DESCRIPTION
     "The total number of connections dropped from the
     backlog by this listener due to all reasons. This
     includes all connections that are allocated initial
     resources, but are not accepted for some reason."
   ::= { tcpEStatsListenerEntry 6 }
tcpEStatsListenerHCSynRcvd OBJECT-TYPE
   SYNTAX ZeroBasedCounter64
   MAX-ACCESS read-only
   STATUS current
   DESCRIPTION
       "The number of SYNs that have been received for this
       listener on systems that can process (or reject) more
       than 1 million connections per second. See
       tcpEStatsListenerSynRcvd."
   ::= { tcpEStatsListenerEntry 7 }
tcpEStatsListenerHCInitial
                            OBJECT-TYPE
  SYNTAX ZeroBasedCounter64
  MAX-ACCESS read-only
  STATUS current
  DESCRIPTION
      "The total number of connections for which the Listener
      has allocated initial state and placed the connection
      in the backlog on systems that can process (or reject)
      more than 1 million connections per second. See
      tcpEStatsListenerInitial."
   ::= { tcpEStatsListenerEntry 8 }
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tcpEStatsListenerHCEstablished OBJECT-TYPE SYNTAX ZeroBasedCounter64 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of connections that have been established to this endpoint on systems that can process (or reject) more than 1 million connections per second. See tcpEStatsListenerEstablished." ::= { tcpEStatsListenerEntry 9 } tcpEStatsListenerHCAccepted OBJECT-TYPE SYNTAX ZeroBasedCounter64 MAX-ACCESS read-only STATUS current DESCRIPTION "The total number of connections for which the Listener has successfully issued an accept, removing the connection from the backlog on systems that can process (or reject) more than 1 million connections per second. See tcpEStatsListenerAccepted." ::= { tcpEStatsListenerEntry 10 } tcpEStatsListenerHCExceedBacklog OBJECT-TYPE SYNTAX ZeroBasedCounter64 MAX-ACCESS read-only STATUS current DESCRIPTION "The total number of connections dropped from the backlog by this listener due to all reasons on systems that can process (or reject) more than 1 million connections per second. See tcpEStatsListenerExceedBacklog." ::= { tcpEStatsListenerEntry 11 } tcpEStatsListenerCurConns OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current DESCRIPTION "The current number of connections in the ESTABLISHED state, which have also been accepted. It excludes connections that have been established but not accepted because they are still subject to being discarded to shed load without explicit action by either endpoint." ::= { tcpEStatsListenerEntry 12 } tcpEStatsListenerMaxBacklog OBJECT-TYPE

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SYNTAX Unsigned32 MAX-ACCESS read-only STATUS current DESCRIPTION "The maximum number of connections allowed in the backlog at one time." ::= { tcpEStatsListenerEntry 13 } tcpEStatsListenerCurBacklog OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current DESCRIPTION "The current number of connections that are in the backlog. This gauge includes connections in ESTABLISHED or SYN-RECEIVED states for which the Listener has not yet issued an accept. If this listener is using some technique to implicitly represent the SYN-RECEIVED states (e.g., by cryptographically encoding the state information in the initial sequence number, ISS), it MAY elect to exclude connections in the SYN-RECEIVED state from the backlog." ::= { tcpEStatsListenerEntry 14 } tcpEStatsListenerCurEstabBacklog OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current DESCRIPTION "The current number of connections in the backlog that are in the ESTABLISHED state, but for which the Listener has not yet issued an accept." ::= { tcpEStatsListenerEntry 15 } _ _ -- TCP Connection ID Table _ _ tcpEStatsConnectIdTable OBJECT-TYPE SYNTAX SEQUENCE OF TcpEStatsConnectIdEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "This table maps information that uniquely identifies each active TCP connection to the connection ID used by

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```
other tables in this MIB Module. It is an extension of
       tcpConnectionTable in RFC 4022.
       Entries are retained in this table for the number of
       seconds indicated by the tcpEStatsConnTableLatency
       object, after the TCP connection first enters the closed
       state."
    ::= { tcpEStats 2 }
tcpEStatsConnectIdEntry OBJECT-TYPE
   SYNTAX TcpEStatsConnectIdEntry
   MAX-ACCESS not-accessible
   STATUS current
   DESCRIPTION
       "Each entry in this table maps a TCP connection
       4-tuple to a connection index."
   AUGMENTS { tcpConnectionEntry }
    ::= { tcpEStatsConnectIdTable 1 }
TcpEStatsConnectIdEntry ::= SEQUENCE {
       tcpEStatsConnectIndex
                                       Unsigned32
}
tcpEStatsConnectIndex OBJECT-TYPE
   SYNTAXUnsigned32 (1..4294967295)MAX-ACCESSread-onlyCTATUCcurrent
   STATUS
                 current
   DESCRIPTION
       "A unique integer value assigned to each TCP Connection
       entry.
       The RECOMMENDED algorithm is to begin at 1 and increase to
       some implementation-specific maximum value and then start
       again at 1 skipping values already in use."
    ::= { tcpEStatsConnectIdEntry 1 }
_ _
-- Basic TCP Performance Statistics
_ _
tcpEStatsPerfTable OBJECT-TYPE
   SYNTAX SEQUENCE OF TcpEStatsPerfEntry
   MAX-ACCESS not-accessible
   STATUS
              current
   DESCRIPTION
       "This table contains objects that are useful for
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measuring TCP performance and first line problem diagnosis. Most objects in this table directly expose some TCP state variable or are easily implemented as simple functions (e.g., the maximum value) of TCP state variables. Entries are retained in this table for the number of seconds indicated by the tcpEStatsConnTableLatency object, after the TCP connection first enters the closed state." ::= { tcpEStats 3 } tcpEStatsPerfEntry OBJECT-TYPE SYNTAX TcpEStatsPerfEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "Each entry in this table has information about the characteristics of each active and recently closed TCP connection." INDEX { tcpEStatsConnectIndex } ::= { tcpEStatsPerfTable 1 } TcpEStatsPerfEntry ::= SEQUENCE { tcpEStatsPerfSegsOutZeroBasedCounter32,tcpEStatsPerfDataSegsOutZeroBasedCounter32,tcpEStatsPerfDataOctetsOutZeroBasedCounter32,tcpEStatsPerfHCDataOctetsOutZeroBasedCounter32,tcpEStatsPerfSegsRetransZeroBasedCounter32,tcpEStatsPerfOctetsRetransZeroBasedCounter32,tcpEStatsPerfSegsInZeroBasedCounter32,tcpEStatsPerfDataOctetsInZeroBasedCounter32,tcpEStatsPerfDataOctetsInZeroBasedCounter32,tcpEStatsPerfElapsedSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfElapsedMicroSecsZeroBasedCounter32,tcpEStatsPerfCurMSSGauge32,tcpEStatsPerfElapsedSizeCounter32, tcpEStatsPerfPipeSizeGauge32,tcpEStatsPerfMaxPipeSizeGauge32,tcpEStatsPerfSmoothedRTTGauge32,Gauge32,Gauge32, tcpEStatsPerfCurRTO Gauge32, tcpEStatsPerfCongSignalsZeroBasedCounter32,tcpEStatsPerfCurCwndGauge32,tcpEStatsPerfCurSsthreshGauge32,tcpEStatsPerfTimeoutsZeroBasedCounter32,tcpEStatsPerfCurRwinSentGauge32,

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tcpEStatsPerfMaxRwinSent Gauge32, tcpEStatsPerfZeroRwinSent ZeroBasedCounter32, tcpEStatsPerfCurRwinRcvd tcpEStatsPerfMaxRwinRcvd Gauge32, Gauge32, tcpEStatsPerfZeroRwinRcvd tcpEStatsPerfZeroRwinRcvd tcpEStatsPerfZeroRwinRcvdZeroBasedCounter32,tcpEStatsPerfSndLimTransRwinZeroBasedCounter32,tcpEStatsPerfSndLimTransCwndZeroBasedCounter32,tcpEStatsPerfSndLimTransSndZeroBasedCounter32,tcpEStatsPerfSndLimTimeRwinZeroBasedCounter32,tcpEStatsPerfSndLimTimeCwndZeroBasedCounter32,tcpEStatsPerfSndLimTimeCwndZeroBasedCounter32,tcpEStatsPerfSndLimTimeSndZeroBasedCounter32, ZeroBasedCounter32, } _ _ -- The following objects provide statistics on aggregate -- segments and data sent on a connection. These provide a -- direct measure of the Internet capacity consumed by a -- connection. _ _ tcpEStatsPerfSegsOut OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only current STATUS DESCRIPTION "The total number of segments sent." ::= { tcpEStatsPerfEntry 1 } tcpEStatsPerfDataSegsOut OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of segments sent containing a positive length data segment." ::= { tcpEStatsPerfEntry 2 } tcpEStatsPerfDataOctetsOut OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "octets" MAX-ACCESS read-only current DESCRIPTION "The number of octets of data contained in transmitted segments, including retransmitted data. Note that this does not include TCP headers." ::= { tcpEStatsPerfEntry 3 }

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tcpEStatsPerfHCDataOctetsOut OBJECT-TYPE SYNTAX ZeroBasedCounter64 "octets" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets of data contained in transmitted segments, including retransmitted data, on systems that can transmit more than 10 million bits per second. Note that this does not include TCP headers." ::= { tcpEStatsPerfEntry 4 } tcpEStatsPerfSegsRetrans OBJECT-TYPE ZeroBasedCounter32 SYNTAX MAX-ACCESS STATUS read-only current DESCRIPTION "The number of segments transmitted containing at least some retransmitted data." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsPerfEntry 5 } tcpEStatsPerfOctetsRetrans OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets retransmitted." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsPerfEntry 6 } tcpEStatsPerfSegsIn OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The total number of segments received." ::= { tcpEStatsPerfEntry 7 } tcpEStatsPerfDataSegsIn OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of segments received containing a positive

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length data segment." ::= { tcpEStatsPerfEntry 8 } tcpEStatsPerfDataOctetsIn OBJECT-TYPE SYNTAX ZeroBasedCounter32 "octets" UNTTS MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets contained in received data segments, including retransmitted data. Note that this does not include TCP headers." ::= { tcpEStatsPerfEntry 9 } tcpEStatsPerfHCDataOctetsIn OBJECT-TYPE SYNTAX ZeroBasedCounter64 "octets" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets contained in received data segments, including retransmitted data, on systems that can receive more than 10 million bits per second. Note that this does not include TCP headers." ::= { tcpEStatsPerfEntry 10 } tcpEStatsPerfElapsedSecs OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "seconds" MAX-ACCESS read-only STATUS current DESCRIPTION "The seconds part of the time elapsed between tcpEStatsPerfStartTimeStamp and the most recent protocol event (segment sent or received)." ::= { tcpEStatsPerfEntry 11 } tcpEStatsPerfElapsedMicroSecs OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "microseconds" MAX-ACCESS read-only STATUS current DESCRIPTION "The micro-second part of time elapsed between tcpEStatsPerfStartTimeStamp to the most recent protocol event (segment sent or received). This may be updated in whatever time granularity is the system supports." ::= { tcpEStatsPerfEntry 12 }

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```
tcpEStatsPerfStartTimeStamp OBJECT-TYPE
   SYNTAX DateAndTime
   MAX-ACCESS
                  read-only
   STATUS
                  current
   DESCRIPTION
      "Time at which this row was created and all
       ZeroBasedCounters in the row were initialized to zero."
    ::= { tcpEStatsPerfEntry 13 }
-- The following objects can be used to fit minimal
  performance models to the TCP data rate.
_ _
_ _
tcpEStatsPerfCurMSS OBJECT-TYPE
   SYNTAX Gauge32
   UNITS
                  "octets"
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
       "The current maximum segment size (MSS), in octets."
   REFERENCE
       "RFC 1122, Requirements for Internet Hosts - Communication
       Layers"
    ::= { tcpEStatsPerfEntry 14 }
tcpEStatsPerfPipeSize OBJECT-TYPE
   SYNTAX Gauge32
UNITS "octets
                  "octets"
   MAX-ACCESS read-only
                  current
   STATUS
   DESCRIPTION
       "The TCP senders current estimate of the number of
       unacknowledged data octets in the network.
       While not in recovery (e.g., while the receiver is not
       reporting missing data to the sender), this is precisely the
       same as 'Flight size' as defined in RFC 2581, which can be
       computed as SND.NXT minus SND.UNA. [RFC793]
       During recovery, the TCP sender has incomplete information
       about the state of the network (e.g., which segments are
       lost vs reordered, especially if the return path is also
       dropping TCP acknowledgments). Current TCP standards do not
       mandate any specific algorithm for estimating the number of
       unacknowledged data octets in the network.
       RFC 3517 describes a conservative algorithm to use SACK
```

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information to estimate the number of unacknowledged data octets in the network. tcpEStatsPerfPipeSize object SHOULD be the same as 'pipe' as defined in RFC 3517 if it is implemented. (Note that while not in recovery the pipe algorithm yields the same values as flight size). If RFC 3517 is not implemented, the data octets in flight SHOULD be estimated as SND.NXT minus SND.UNA adjusted by some measure of the data that has left the network and retransmitted data. For example, with Reno or NewReno style TCP, the number of duplicate acknowledgment is used to count the number of segments that have left the network. That is, PipeSize=SND.NXT-SND.UNA+(retransmits-dupacks)*CurMSS" REFERENCE "RFC 793, RFC 2581, RFC 3517" ::= { tcpEStatsPerfEntry 15 } tcpEStatsPerfMaxPipeSize OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" read-only MAX-ACCESS STATUS current DESCRIPTION "The maximum value of tcpEStatsPerfPipeSize, for this connection." REFERENCE "RFC 793, RFC 2581, RFC 3517" ::= { tcpEStatsPerfEntry 16 } tcpEStatsPerfSmoothedRTT OBJECT-TYPE SYNTAX Gauge32 UNITS "milliseconds" MAX-ACCESS read-only STATUS current DESCRIPTION "The smoothed round trip time used in calculation of the RTO. See SRTT in [RFC2988]." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPerfEntry 17 } tcpEStatsPerfCurRTO OBJECT-TYPE SYNTAX Gauge32 "milliseconds" UNITS

UNITS "milliseconds MAX-ACCESS read-only STATUS current DESCRIPTION

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"The current value of the retransmit timer RTO." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPerfEntry 18 } tcpEStatsPerfCongSignals OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only SYNTAA MAX-ACCESS STATUS current DESCRIPTION "The number of multiplicative downward congestion window adjustments due to all forms of congestion signals, including Fast Retransmit, Explicit Congestion Notification (ECN), and timeouts. This object summarizes all events that invoke the MD portion of Additive Increase Multiplicative Decrease (AIMD) congestion control, and as such is the best indicator of how a cwnd is being affected by congestion. Note that retransmission timeouts multiplicatively reduce the window implicitly by setting ssthresh, and SHOULD be included in tcpEStatsPerfCongSignals. In order to minimize spurious congestion indications due to out-of-order segments, tcpEStatsPerfCongSignals SHOULD be incremented in association with the Fast Retransmit algorithm." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPerfEntry 19 } tcpEStatsPerfCurCwnd OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The current congestion window, in octets." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPerfEntry 20 } tcpEStatsPerfCurSsthresh OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The current slow start threshold in octets." REFERENCE "RFC 2581, TCP Congestion Control" Mathis, et al. Standards Track [Page 24]

::= { tcpEStatsPerfEntry 21 } tcpEStatsPerfTimeouts OBJECT-TYPE ZeroBasedCounter32 SYNIAA MAX-ACCESS SYNTAX read-only STATUS current DESCRIPTION "The number of times the retransmit timeout has expired when the RTO backoff multiplier is equal to one." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPerfEntry 22 } -- The following objects instrument receiver window updates -- sent by the local receiver to the remote sender. These can -- be used to determine if the local receiver is exerting flow -- control back pressure on the remote sender. _ _ tcpEStatsPerfCurRwinSent OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" read-only MAX-ACCESS STATUS current DESCRIPTION "The most recent window advertisement sent, in octets." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsPerfEntry 23 } tcpEStatsPerfMaxRwinSent OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The maximum window advertisement sent, in octets." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsPerfEntry 24 } tcpEStatsPerfZeroRwinSent OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of acknowledgments sent announcing a zero

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```
receive window, when the previously announced window was
       not zero."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
   ::= { tcpEStatsPerfEntry 25 }
-- The following objects instrument receiver window updates
-- from the far end-system to determine if the remote receiver
-- has sufficient buffer space or is exerting flow-control
-- back pressure on the local sender.
_ _
tcpEStatsPerfCurRwinRcvd OBJECT-TYPE
   SYNTAX Gauge32
   UNITS
                  "octets"
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
      "The most recent window advertisement received, in octets."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
   ::= { tcpEStatsPerfEntry 26 }
tcpEStatsPerfMaxRwinRcvd OBJECT-TYPE
   SYNTAX Gauge32
UNITS "octets"
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
      "The maximum window advertisement received, in octets."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
   ::= { tcpEStatsPerfEntry 27 }
tcpEStatsPerfZeroRwinRcvd OBJECT-TYPE
   SYNTAX ZeroBasedCounter32
   MAX-ACCESS read-only
   STATUS
                 current
   DESCRIPTION
      "The number of acknowledgments received announcing a zero
       receive window, when the previously announced window was
       not zero."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
   ::= { tcpEStatsPerfEntry 28 }
```

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The following optional objects can be used to quickly -- identify which subsystems are limiting TCP performance. -- There are three parallel pairs of instruments that measure -- the extent to which TCP performance is limited by the -- announced receiver window (indicating a receiver -- bottleneck), the current congestion window or -- retransmission timeout (indicating a path bottleneck) and -- all others events (indicating a sender bottleneck). _ _ -- These instruments SHOULD be updated every time the TCP -- output routine stops sending data. The elapsed time since -- the previous stop is accumulated into the appropriate -- object as determined by the previous stop reason (e.g., -- stop state). The current stop reason determines which timer -- will be updated the next time TCP output stops. _ _ -- Since there is no explicit stop at the beginning of a -- timeout, it is necessary to retroactively reclassify the -- previous stop as 'Congestion Limited'. _ _ tcpEStatsPerfSndLimTransRwin OBJECT-TYPE SYNTAXZeroBasedCounter32MAX-ACCESSread-only STATUS current DESCRIPTION "The number of transitions into the 'Receiver Limited' state from either the 'Congestion Limited' or 'Sender Limited' states. This state is entered whenever TCP transmission stops because the sender has filled the announced receiver window, i.e., when SND.NXT has advanced to SND.UNA + SND.WND - 1 as described in RFC 793." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsPerfEntry 31 } tcpEStatsPerfSndLimTransCwnd OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of transitions into the 'Congestion Limited' state from either the 'Receiver Limited' or 'Sender Limited' states. This state is entered whenever TCP transmission stops because the sender has reached some limit defined by congestion control (e.g., cwnd) or other algorithms (retransmission timeouts) designed to control network traffic. See the definition of 'CONGESTION WINDOW'

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in RFC 2581."
REFERENCE
 "RFC 2581, TCP Congestion Control"
 ::= { tcpEStatsPerfEntry 32 }
tcpEStatsPerfSndLimTransSnd OBJECT-TYPE
 SYNTAX ZeroBasedCounter32
 MAX-ACCESS read-only
 STATUS current
 DESCRIPTION
 "The number of transitions into the 'Sender Limited' state
 from either the 'Receiver Limited' or 'Congestion Limited'

from either the 'Receiver Limited' or 'Congestion Limited'
states. This state is entered whenever TCP transmission
stops due to some sender limit such as running out of
application data or other resources and the Karn algorithm.
When TCP stops sending data for any reason, which cannot be
classified as Receiver Limited or Congestion Limited, it
MUST be treated as Sender Limited."
::= { tcpEStatsPerfEntry 33 }

tcpEStatsPerfSndLimTimeRwin OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "milliseconds" MAX-ACCESS read-only STATUS current DESCRIPTION "The cumulative time spent in the 'Receiver Limited' state. See tcpEStatsPerfSndLimTransRwin." ::= { tcpEStatsPerfEntry 34 }

```
tcpEStatsPerfSndLimTimeCwnd OBJECT-TYPE
SYNTAX ZeroBasedCounter32
UNITS "milliseconds"
MAX-ACCESS read-only
STATUS current
DESCRIPTION
"The cumulative time spent in the 'Congestion Limited'
state. See tcpEStatsPerfSndLimTransCwnd. When there is a
retransmission timeout, it SHOULD be counted in
tcpEStatsPerfSndLimTimeCwnd (and not the cumulative time
for some other state.)"
::= { tcpEStatsPerfEntry 35 }
```

tcpEStatsPerfSndLimTimeSnd OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "milliseconds" MAX-ACCESS read-only STATUS current

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```
DESCRIPTION
       "The cumulative time spent in the 'Sender Limited' state.
        See tcpEStatsPerfSndLimTransSnd."
    ::= { tcpEStatsPerfEntry 36 }
-- Statistics for diagnosing path problems
_ _
tcpEStatsPathTable OBJECT-TYPE
    SYNTAX SEQUENCE OF TCPEStatsPathEntry
    MAX-ACCESS not-accessible
    STATUS current
    DESCRIPTION
        "This table contains objects that can be used to infer
        detailed behavior of the Internet path, such as the
        extent that there is reordering, \ensuremath{\mathsf{ECN}} bits, and if
        RTT fluctuations are correlated to losses.
        Entries are retained in this table for the number of
        seconds indicated by the tcpEStatsConnTableLatency
        object, after the TCP connection first enters the closed
        state."
    ::= { tcpEStats 4 }
tcpEStatsPathEntry OBJECT-TYPE
    SYNTAX TcpEStatsPathEntry
MAX-ACCESS not-accessible
    STATUS current
    DESCRIPTION
        "Each entry in this table has information about the
        characteristics of each active and recently closed TCP
        connection."
   INDEX { tcpEStatsConnectIndex }
   ::= { tcpEStatsPathTable 1 }
TcpEStatsPathEntry ::= SEQUENCE {
        COPESCALSPATNRetranThreshGauge32,tcpEStatsPathNonRecovDAEpisodesZeroBasedCounter32,tcpEStatsPathSumOctetsReorderedZeroBasedCounter32,tcpEStatsPathNonRecovDAZeroBasedCounter32,
        tcpEStatsPathSampleRTT
                                              Gauge32,
        tcpEStatsPathRTTVar
                                              Gauge32,
        tcpEStatsPathMaxRTT
                                              Gauge32,
        tcpEStatsPathMinRTT
                                              Gauge32,
        tcpEStatsPathSumRTT
                                              ZeroBasedCounter32,
```

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tcpEStatsPathHCSumRTT ZeroBasedCounter64, tcpEStatsPathCountRTT ZeroBasedCounter32, Gauge32, tcpEStatsPathMaxRTO Gauge32, tcpEStatsPathMinRTO Unsigned32, tcpEStatsPathIpTtl tcpEStatsPathIpTosIn tcpEStatsPathIpTosOut OCTET STRING, tcpEStatsPathIpTosOutOCTET STRING,tcpEStatsPathIpTosOutOCTET STRING,tcpEStatsPathPreCongSumCwndZeroBasedCounter32,tcpEStatsPathPreCongSumRTTZeroBasedCounter32,tcpEStatsPathPostCongSumRTTZeroBasedCounter32,tcpEStatsPathPostCongCountRTTZeroBasedCounter32,tcpEStatsPathECNsignalsZeroBasedCounter32,tcpEStatsPathDupAckEpisodesZeroBasedCounter32,tcpEStatsPathDupAckEpisodesZeroBasedCounter32, tcpEStatsPathDupAckEpisouestcpEStatsPathRcvRTTGauge32,tcpEStatsPathDupAcksOutZeroBasedCounter32,tcpEStatsPathCERcvdZeroBasedCounter32,TototsPathECESentZeroBasedCounter32 } _ _ -- The following optional objects can be used to infer segment reordering on the path from the local sender to the remote _ _ receiver. tcpEStatsPathRetranThresh OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current current STATUS DESCRIPTION "The number of duplicate acknowledgments required to trigger Fast Retransmit. Note that although this is constant in traditional Reno TCP implementations, it is adaptive in many newer TCPs." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPathEntry 1 } tcpEStatsPathNonRecovDAEpisodes OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of duplicate acknowledgment episodes that did not trigger a Fast Retransmit because ACK advanced prior to the number of duplicate acknowledgments reaching RetranThresh.

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In many implementations this is the number of times the 'dupacks' counter is set to zero when it is non-zero but less than RetranThresh. Note that the change in tcpEStatsPathNonRecovDAEpisodes divided by the change in tcpEStatsPerfDataSegsOut is an estimate of the frequency of data reordering on the forward path over some interval." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPathEntry 2 } tcpEStatsPathSumOctetsReordered OBJECT-TYPE SYNTAX ZeroBasedCounter32 "octets" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The sum of the amounts SND.UNA advances on the acknowledgment which ends a dup-ack episode without a retransmission. Note the change in tcpEStatsPathSumOctetsReordered divided by the change in tcpEStatsPathNonRecovDAEpisodes is an estimates of the average reordering distance, over some interval." ::= { tcpEStatsPathEntry 3 } tcpEStatsPathNonRecovDA OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only current STATUS DESCRIPTION "Duplicate acks (or SACKS) that did not trigger a Fast Retransmit because ACK advanced prior to the number of duplicate acknowledgments reaching RetranThresh. In many implementations, this is the sum of the 'dupacks' counter, just before it is set to zero because ACK advanced without a Fast Retransmit. Note that the change in tcpEStatsPathNonRecovDA divided by the change in tcpEStatsPathNonRecovDAEpisodes is an estimate of the average reordering distance in segments over some interval." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPathEntry 4 }

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_ _

```
-- The following optional objects instrument the round trip
-- time estimator and the retransmission timeout timer.
_ _
tcpEStatsPathSampleRTT OBJECT-TYPE
   SYNTAX Gauge32
UNITS "milliseconds"
   MAX-ACCESS
                 read-only
   STATUS
                  current
   DESCRIPTION
      "The most recent raw round trip time measurement used in
       calculation of the RTO."
   REFERENCE
      "RFC 2988, Computing TCP's Retransmission Timer"
    ::= { tcpEStatsPathEntry 11 }
tcpEStatsPathRTTVar OBJECT-TYPE
   SYNTAX Gauge32
                  "milliseconds"
   UNITS
   MAX-ACCESS read-only
STATUS
                  current
   STATUS
   DESCRIPTION
       "The round trip time variation used in calculation of the
       RTO. See RTTVAR in [RFC2988]."
   REFERENCE
      "RFC 2988, Computing TCP's Retransmission Timer"
    ::= { tcpEStatsPathEntry 12 }
tcpEStatsPathMaxRTT OBJECT-TYPE
   SYNTAX Gauge32
UNITS "milliseconds"
   MAX-ACCESS read-only
   STATUS
                 current
   DESCRIPTION
      "The maximum sampled round trip time."
   REFERENCE
      "RFC 2988, Computing TCP's Retransmission Timer"
    ::= { tcpEStatsPathEntry 13 }
tcpEStatsPathMinRTT OBJECT-TYPE
   SYNTAX Gauge32
   UNITS
                  "milliseconds"
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
      "The minimum sampled round trip time."
   REFERENCE
```

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"RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 14 } tcpEStatsPathSumRTT OBJECT-TYPE SYNTAXZeroBasedCounter32UNITS"milliseconds"MAX-ACCESSread-onlySTATUScurrent current DESCRIPTION "The sum of all sampled round trip times. Note that the change in tcpEStatsPathSumRTT divided by the change in tcpEStatsPathCountRTT is the mean RTT, uniformly averaged over an enter interval." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 15 } tcpEStatsPathHCSumRTT OBJECT-TYPE SYNTAX ZeroBasedCounter64 "milliseconds" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The sum of all sampled round trip times, on all systems that implement multiple concurrent RTT measurements. Note that the change in tcpEStatsPathHCSumRTT divided by the change in tcpEStatsPathCountRTT is the mean RTT, uniformly averaged over an enter interval." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 16 } tcpEStatsPathCountRTT OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of round trip time samples included in tcpEStatsPathSumRTT and tcpEStatsPathHCSumRTT." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 17 } tcpEStatsPathMaxRTO OBJECT-TYPE SYNTAX Gauge32 UNITS "milliseconds"

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MAX-ACCESS read-only STATUS current DESCRIPTION "The maximum value of the retransmit timer RTO." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 18 } tcpEStatsPathMinRTO OBJECT-TYPE SYNTAX Gauge32 UNITS "milliseconds" MAX-ACCESS read-only STATUS current DESCRIPTION "The minimum value of the retransmit timer RTO." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsPathEntry 19 } _ _ -- The following optional objects provide information about how TCP is using the IP layer. ___ tcpEStatsPathIpTtl OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS read-only STATUS current DESCRIPTION "The value of the TTL field carried in the most recently received IP header. This is sometimes useful to detect changing or unstable routes." REFERENCE "RFC 791, Internet Protocol" ::= { tcpEStatsPathEntry 20 } tcpEStatsPathIpTosIn OBJECT-TYPE SYNTAX OCTET STRING (SIZE(1)) MAX-ACCESS read-only STATUS current DESCRIPTION "The value of the IPv4 Type of Service octet, or the IPv6 traffic class octet, carried in the most recently received IP header. This is useful to diagnose interactions between TCP and any IP layer packet scheduling and delivery policy, which might be in effect to implement Diffserv."

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REFERENCE "RFC 3260, New Terminology and Clarifications for Diffserv" ::= { tcpEStatsPathEntry 21 } tcpEStatsPathIpTosOut OBJECT-TYPE SYNTAX OCTET STRING (SIZE(1)) MAX-ACCESS STATUS read-only current DESCRIPTION "The value of the IPv4 Type Of Service octet, or the IPv6 traffic class octet, carried in the most recently transmitted IP header. This is useful to diagnose interactions between TCP and any IP layer packet scheduling and delivery policy, which might be in effect to implement Diffserv." REFERENCE "RFC 3260, New Terminology and Clarifications for Diffserv" ::= { tcpEStatsPathEntry 22 } _ _ -- The following optional objects characterize the congestion feedback signals by collecting statistics on how the _ _ congestion events are correlated to losses, changes in RTT _ _ -- and other protocol events. _ _ tcpEStatsPathPreCongSumCwnd OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "octets" UNITS read-only STATUS current DESCRIPTION "The sum of the values of the congestion window, in octets, captured each time a congestion signal is received. This MUST be updated each time tcpEStatsPerfCongSignals is incremented, such that the change in tcpEStatsPathPreCongSumCwnd divided by the change in tcpEStatsPerfCongSignals is the average window (over some interval) just prior to a congestion signal." ::= { tcpEStatsPathEntry 23 } tcpEStatsPathPreCongSumRTT OBJECT-TYPE SYNTAX ZeroBasedCounter32 "milliseconds" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION

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"Sum of the last sample of the RTT (tcpEStatsPathSampleRTT) prior to the received congestion signals. This MUST be updated each time tcpEStatsPerfCongSignals is incremented, such that the change in tcpEStatsPathPreCongSumRTT divided by the change in tcpEStatsPerfCongSignals is the average RTT (over some interval) just prior to a congestion signal." ::= { tcpEStatsPathEntry 24 } tcpEStatsPathPostCongSumRTT OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "Sum of the first sample of the RTT (tcpEStatsPathSampleRTT) following each congestion signal. Such that the change in tcpEStatsPathPostCongSumRTT divided by the change in tcpEStatsPathPostCongCountRTT is the average RTT (over some interval) just after a congestion signal." ::= { tcpEStatsPathEntry 25 } tcpEStatsPathPostCongCountRTT OBJECT-TYPE SYNTAXZeroBasedCounter32UNITS"milliseconds"MAX-ACCESSread-onlySTATUScurrent STATUS current DESCRIPTION "The number of RTT samples included in tcpEStatsPathPostCongSumRTT such that the change in tcpEStatsPathPostCongSumRTT divided by the change in tcpEStatsPathPostCongCountRTT is the average RTT (over some interval) just after a congestion signal." ::= { tcpEStatsPathEntry 26 } -- The following optional objects can be used to detect other -- types of non-loss congestion signals such as source quench -- or ECN. _ _ tcpEStatsPathECNsignals OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of congestion signals delivered to the TCP sender via explicit congestion notification (ECN). This is typically the number of segments bearing Echo Congestion

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Experienced (ECE) bits, but should also include segments failing the ECN nonce check or other explicit congestion signals." REFERENCE "RFC 3168, The Addition of Explicit Congestion Notification (ECN) to IP" ::= { tcpEStatsPathEntry 27 } -- The following optional objects are receiver side -- instruments of the path from the sender to the receiver. In -- general, the receiver has less information about the state -- of the path because the receiver does not have a robust -- mechanism to infer the sender's actions. _ _ tcpEStatsPathDupAckEpisodes OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only current STATUS DESCRIPTION "The number of Duplicate Acks Sent when prior Ack was not duplicate. This is the number of times that a contiguous series of duplicate acknowledgments have been sent. This is an indication of the number of data segments lost or reordered on the path from the remote TCP endpoint to the near TCP endpoint." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPathEntry 28 } tcpEStatsPathRcvRTT OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current DESCRIPTION "The receiver's estimate of the Path RTT. Adaptive receiver window algorithms depend on the receiver to having a good estimate of the path RTT." ::= { tcpEStatsPathEntry 29 } tcpEStatsPathDupAcksOut OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION

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"The number of duplicate ACKs sent. The ratio of the change in tcpEStatsPathDupAcksOut to the change in tcpEStatsPathDupAckEpisodes is an indication of reorder or recovery distance over some interval." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsPathEntry 30 } tcpEStatsPathCERcvd OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS STATUS read-only current DESCRIPTION "The number of segments received with IP headers bearing Congestion Experienced (CE) markings." REFERENCE "RFC 3168, The Addition of Explicit Congestion Notification (ECN) to IP" ::= { tcpEStatsPathEntry 31 } tcpEStatsPathECESent OBJECT-TYPE SYNTAX ZeroBasedCounter32 read-only MAX-ACCESS current STATUS DESCRIPTION "Number of times the Echo Congestion Experienced (ECE) bit in the TCP header has been set (transitioned from 0 to 1), due to a Congestion Experienced (CE) marking on an IP header. Note that ECE can be set and reset only once per RTT, while CE can be set on many segments per RTT." REFERENCE "RFC 3168, The Addition of Explicit Congestion Notification (ECN) to IP" ::= { tcpEStatsPathEntry 32 } _ _ -- Statistics for diagnosing stack algorithms _ _ tcpEStatsStackTable OBJECT-TYPE SYNTAX SEQUENCE OF TcpEStatsStackEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "This table contains objects that are most useful for determining how well some of the TCP control algorithms are coping with this particular

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```
path.
            Entries are retained in this table for the number of
            seconds indicated by the tcpEStatsConnTableLatency
            object, after the TCP connection first enters the closed
            state."
      ::= { tcpEStats 5 }
tcpEStatsStackEntry OBJECT-TYPE
      SYNTAX TcpEStatsStackEntry
      MAX-ACCESS not-accessible
      STATUS current
      DESCRIPTION
            "Each entry in this table has information about the
            characteristics of each active and recently closed TCP
            connection."
     INDEX { tcpEStatsConnectIndex }
     ::= { tcpEStatsStackTable 1 }
TcpEStatsStackEntry ::= SEQUENCE {
            tcpEStatsStackActiveOpen
                                                                  TruthValue,
            tcpEStatsStackMSSSent
                                                                   Unsigned32,
            tcpEStatsStackMSSRcvdUnsigned32,tcpEStatsStackWinScaleSentInteger32,tcpEStatsStackWinScaleRcvdInteger32,tcpEStatsStackTimeStampsTcpEStatsNegotiated,tcpEStatsStackECNTcpEStatsNegotiated,
                                                                 TcpEStatsNegotiated,
            tcpEStatsStackECN
            tcpEStatsStackWillSendSACK TcpEStatsNegotiated,
tcpEStatsStackWillUseSACK TcpEStatsNegotiated,
tcpEStatsStackWillUseSACK TcpEStatsNegotiated,
                                                                  INTEGER,
            tcpEStatsStackState
                                                                  TruthValue,
            tcpEStatsStackNagle
            tcpEStatsStackMaxSsCwnd
tcpEStatsStackMaxCaCwnd
                                                                 Gauge32,
                                                                 Gauge32,
            tcpEStatsStackMaxCacwindGauge32,tcpEStatsStackMaxSsthreshGauge32,tcpEStatsStackMinSsthreshGauge32,tcpEStatsStackInRecoveryINTEGER,tcpEStatsStackDupAcksInZeroBased
            tcpEStatsStackDupAcksIn ZeroBasedCounter32,
tcpEStatsStackSpuriousFrDetected ZeroBasedCounter32,
            tcpEStatsStackSpuriousRtoDetected ZeroBasedCounter32,
            tcpEStatsStackSoftErrorsZeroBasedCounter32,tcpEStatsStackSoftErrorReasonINTEGER,tcpEStatsStackSlowStartZeroBasedCounter32,tcpEStatsStackCongAvoidZeroBasedCounter32,
            tcpEStatsStackCongAvoidZeroBasedCounter32,tcpEStatsStackOtherReductionsZeroBasedCounter32,tcpEStatsStackCongOverCountZeroBasedCounter32,tcpEStatsStackFastRetranZeroBasedCounter32,tcpEStatsStackSubsequentTimeouteZeroBasedCounter32,
            tcpEStatsStackSubsequentTimeouts ZeroBasedCounter32,
```

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<pre>tcpEStatsStackCurTimeoutCount tcpEStatsStackAbruptTimeouts tcpEStatsStackSACKsRcvd tcpEStatsStackSACKBlocksRcvd tcpEStatsStackSendStall tcpEStatsStackDSACKDups tcpEStatsStackMaxMSS tcpEStatsStackMinMSS tcpEStatsStackSndInitial tcpEStatsStackRecInitial tcpEStatsStackCurRetxQueue tcpEStatsStackCurRetxQueue tcpEStatsStackCurReasmQueue tcpEStatsStackMaxReasmQueue }</pre>	Gauge32, ZeroBasedCounter32, ZeroBasedCounter32, ZeroBasedCounter32, ZeroBasedCounter32, Gauge32, Gauge32, Unsigned32, Unsigned32, Gauge32, Gauge32, Gauge32, Gauge32, Gauge32, Gauge32, Gauge32,			
 The following objects reflect TCP options carried on the SYN or SYN-ACK. These options are used to provide additional protocol parameters or to enable various optional TCP features or algorithms. Except as noted, the TCP protocol does not permit these options to change after the SYN exchange. 				
 tcpEStatsStackActiveOpen OBJECT-TYPE SYNTAX TruthValue MAX-ACCESS read-only STATUS current DESCRIPTION "True(1) if the local connection state, else false(2)." REFERENCE "RFC 793, Transmission Control Pr ::= { tcpEStatsStackEntry 1 }				
<pre>tcpEStatsStackMSSSent OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS read-only STATUS current DESCRIPTION "The value sent in an MSS option, REFERENCE "RFC 1122, Requirements for Inter Layers" ::= { tcpEStatsStackEntry 2 }</pre>				

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tcpEStatsStackMSSRcvd OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS read-only STATUS current DESCRIPTION "The value received in an MSS option, or zero if none." REFERENCE "RFC 1122, Requirements for Internet Hosts - Communication Lavers" ::= { tcpEStatsStackEntry 3 } tcpEStatsStackWinScaleSent OBJECT-TYPE SYNTAX Integer32 (-1..14) MAX-ACCESS read-only STATUS current DESCRIPTION "The value of the transmitted window scale option if one was sent; otherwise, a value of -1. Note that if both tcpEStatsStackWinScaleSent and tcpEStatsStackWinScaleRcvd are not -1, then Rcv.Wind.Scale will be the same as this value and used to scale receiver window announcements from the local host to the remote host." REFERENCE "RFC 1323, TCP Extensions for High Performance" ::= { tcpEStatsStackEntry 4 } tcpEStatsStackWinScaleRcvd OBJECT-TYPE SYNTAXInteger32 (-1..14)MAX-ACCESSread-only current STATUS DESCRIPTION "The value of the received window scale option if one was received; otherwise, a value of -1. Note that if both tcpEStatsStackWinScaleSent and tcpEStatsStackWinScaleRcvd are not -1, then Snd.Wind.Scale will be the same as this value and used to scale receiver window announcements from the remote host to the local host." REFERENCE "RFC 1323, TCP Extensions for High Performance" ::= { tcpEStatsStackEntry 5 } tcpEStatsStackTimeStamps OBJECT-TYPE SYNTAX TcpEStatsNegotiated MAX-ACCESS read-only Mathis, et al. Standards Track [Page 41]

STATUS current DESCRIPTION "Enabled(1) if TCP timestamps have been negotiated on, selfDisabled(2) if they are disabled or not implemented on the local host, or peerDisabled(3) if not negotiated by the remote hosts." REFERENCE "RFC 1323, TCP Extensions for High Performance" ::= { tcpEStatsStackEntry 6 } tcpEStatsStackECN OBJECT-TYPE TcpEStatsNegotiated read-only SYNTAX MAX-ACCESS STATUS current DESCRIPTION "Enabled(1) if Explicit Congestion Notification (ECN) has been negotiated on, selfDisabled(2) if it is disabled or not implemented on the local host, or peerDisabled(3) if not negotiated by the remote hosts." REFERENCE "RFC 3168, The Addition of Explicit Congestion Notification (ECN) to IP" ::= { tcpEStatsStackEntry 7 } tcpEStatsStackWillSendSACK OBJECT-TYPE SYNTAX TcpEStatsNegotiated read-only MAX-ACCESS STATUS current DESCRIPTION "Enabled(1) if the local host will send SACK options, selfDisabled(2) if SACK is disabled or not implemented on the local host, or peerDisabled(3) if the remote host did not send the SACK-permitted option. Note that SACK negotiation is not symmetrical. SACK can enabled on one side of the connection and not the other." REFERENCE "RFC 2018, TCP Selective Acknowledgement Options" ::= { tcpEStatsStackEntry 8 } tcpEStatsStackWillUseSACK OBJECT-TYPE SYNTAX TcpEStatsNegotiated MAX-ACCESS read-only STATUS current DESCRIPTION "Enabled(1) if the local host will process SACK options, selfDisabled(2) if SACK is disabled or not implemented on the local host, or peerDisabled(3) if the remote host sends

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duplicate ACKs without SACK options, or the local host otherwise decides not to process received SACK options.

Unlike other TCP options, the remote data receiver cannot explicitly indicate if it is able to generate SACK options. When sending data, the local host has to deduce if the remote receiver is sending SACK options. This object can transition from Enabled(1) to peerDisabled(3) after the SYN exchange.

Note that SACK negotiation is not symmetrical. SACK can enabled on one side of the connection and not the other." REFERENCE

"RFC 2018, TCP Selective Acknowledgement Options"
::= { tcpEStatsStackEntry 9 }

-- The following two objects reflect the current state of the -- connection.

```
_ _
```

_ _

tcpEStatsStackState OBJECT-TYPE SYNTAX INTEGER { tcpESStateClosed(1), tcpESStateListen(2), tcpESStateSynSent(3), tcpESStateSynReceived(4), tcpESStateEstablished(5), tcpESStateFinWait1(6), tcpESStateFinWait2(7), tcpESStateCloseWait(8), tcpESStateLastAck(9), tcpESStateClosing(10), tcpESStateTimeWait(11), tcpESStateDeleteTcb(12) } MAX-ACCESS read-only STATUS current DESCRIPTION "An integer value representing the connection state from the TCP State Transition Diagram. The value listen(2) is included only for parallelism to the old tcpConnTable, and SHOULD NOT be used because the listen state in managed by the tcpListenerTable. The value DeleteTcb(12) is included only for parallelism to the tcpConnTable mechanism for terminating connections,

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```
although this table does not permit writing."
   REFERENCE
     "RFC 793, Transmission Control Protocol"
    ::= { tcpEStatsStackEntry 10 }
tcpEStatsStackNagle OBJECT-TYPE
   SYNTAX TruthValue
MAX-ACCESS read-only
   MAX-ACCESS
   STATUS
                  current
   DESCRIPTION
       "True(1) if the Nagle algorithm is being used, else
       false(2)."
   REFERENCE
      "RFC 1122, Requirements for Internet Hosts - Communication
       Layers"
   ::= { tcpEStatsStackEntry 11 }
_ _
-- The following objects instrument the overall operation of
  TCP congestion control and data retransmissions. These
_ _
-- instruments are sufficient to fit the actual performance to
   an updated macroscopic performance model [RFC2581] [Mat97]
_ _
   [Pad98].
_ _
_ _
tcpEStatsStackMaxSsCwnd OBJECT-TYPE
   SYNTAX Gauge32
UNITS "octets"
   MAX-ACCESS read-only
   STATUS
                   current
   DESCRIPTION
      "The maximum congestion window used during Slow Start, in
       octets."
   REFERENCE
      "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 12 }
tcpEStatsStackMaxCaCwnd OBJECT-TYPE
   SYNTAX Gauge32
   UNITS
                   "octets"
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
      "The maximum congestion window used during Congestion
       Avoidance, in octets."
   REFERENCE
      "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 13 }
```

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tcpEStatsStackMaxSsthresh OBJECT-TYPE SYNTAX Gauge32 "octets" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The maximum slow start threshold, excluding the initial value." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsStackEntry 14 } tcpEStatsStackMinSsthresh OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The minimum slow start threshold." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsStackEntry 15 } tcpEStatsStackInRecovery OBJECT-TYPE SYNTAX INTEGER { tcpESDataContiguous(1), tcpESDataUnordered(2), tcpESDataRecovery(3) } MAX-ACCESS read-only STATUS current DESCRIPTION "An integer value representing the state of the loss recovery for this connection. tcpESDataContiguous(1) indicates that the remote receiver is reporting contiguous data (no duplicate acknowledgments or SACK options) and that there are no unacknowledged retransmissions. tcpESDataUnordered(2) indicates that the remote receiver is reporting missing or out-of-order data (e.g., sending duplicate acknowledgments or SACK options) and that there are no unacknowledged retransmissions (because the missing data has not yet been retransmitted). tcpESDataRecovery(3) indicates that the sender has outstanding retransmitted data that is still

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```
unacknowledged."
    REFERENCE
      "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 16 }
tcpEStatsStackDupAcksIn OBJECT-TYPE
   SYNTAXZeroBasedCounter32MAX-ACCESSread-onlySTATUScurrent
   STATUS
                   current
    DESCRIPTION
       "The number of duplicate ACKs received."
    REFERENCE
       "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 17 }
tcpEStatsStackSpuriousFrDetected OBJECT-TYPE
   SYNTAXZeroBasedCounter32MAX-ACCESSread-onlySTATUScurrent
                   current
   STATUS
   DESCRIPTION
       "The number of acknowledgments reporting out-of-order
        segments after the Fast Retransmit algorithm has already
        retransmitted the segments. (For example as detected by the
        Eifel algorithm).'"
    REFERENCE
       "RFC 3522, The Eifel Detection Algorithm for TCP"
    ::= { tcpEStatsStackEntry 18 }
tcpEStatsStackSpuriousRtoDetected OBJECT-TYPE
   SYNTAXZeroBasedCounter32MAX-ACCESSread-only
                   current
   STATUS
   DESCRIPTION
       "The number of acknowledgments reporting segments that have
       already been retransmitted due to a Retransmission Timeout."
    ::= { tcpEStatsStackEntry 19 }
_ _
-- The following optional objects instrument unusual protocol
-- events that probably indicate implementation problems in
-- the protocol or path.
_ _
tcpEStatsStackSoftErrors OBJECT-TYPE
   SYNTAX ZeroBasedCounter32
   MAX-ACCESS read-only
    STATUS
                   current
   DESCRIPTION
```

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```
"The number of segments that fail various consistency tests
       during TCP input processing. Soft errors might cause the
       segment to be discarded but some do not. Some of these soft
       errors cause the generation of a TCP acknowledgment, while
       others are silently discarded."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
    ::= { tcpEStatsStackEntry 21 }
tcpEStatsStackSoftErrorReason OBJECT-TYPE
   SYNTAX
                  INTEGER {
      belowDataWindow(1),
      aboveDataWindow(2),
      belowAckWindow(3),
      aboveAckWindow(4),
      belowTSWindow(5),
      aboveTSWindow(6),
      dataCheckSum(7),
      otherSoftError(8)
   }
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
       "This object identifies which consistency test most recently
       failed during TCP input processing. This object SHOULD be
       set every time tcpEStatsStackSoftErrors is incremented. The
       codes are as follows:
       belowDataWindow(1) - All data in the segment is below
       SND.UNA. (Normal for keep-alives and zero window probes).
       aboveDataWindow(2) - Some data in the segment is above
       SND.WND. (Indicates an implementation bug or possible
       attack).
       belowAckWindow(3) - ACK below SND.UNA. (Indicates that the
       return path is reordering ACKs)
       aboveAckWindow(4) - An ACK for data that we have not sent.
       (Indicates an implementation bug or possible attack).
       belowTSWindow(5) - TSecr on the segment is older than the
       current TS.Recent (Normal for the rare case where PAWS
       detects data reordered by the network).
       aboveTSWindow(6) - TSecr on the segment is newer than the
       current TS.Recent. (Indicates an implementation bug or
       possible attack).
```

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```
dataCheckSum(7) - Incorrect checksum. Note that this value
        is intrinsically fragile, because the header fields used to
        identify the connection may have been corrupted.
       otherSoftError(8) - All other soft errors not listed
       above."
   REFERENCE
      "RFC 793, Transmission Control Protocol"
    ::= { tcpEStatsStackEntry 22 }
-- The following optional objects expose the detailed
-- operation of the congestion control algorithms.
_ _
tcpEStatsStackSlowStart OBJECT-TYPE
   SYNTAX ZeroBasedCounter32
                  read-only
   MAX-ACCESS
                  current
   STATUS
   DESCRIPTION
       "The number of times the congestion window has been
       increased by the Slow Start algorithm."
   REFERENCE
       "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 23 }
tcpEStatsStackCongAvoid OBJECT-TYPE
   SYNTAXZeroBasedCounter32MAX-ACCESSread-onlySTATUSSuprement
                  current
   STATUS
   DESCRIPTION
       "The number of times the congestion window has been
       increased by the Congestion Avoidance algorithm."
   REFERENCE
       "RFC 2581, TCP Congestion Control"
    ::= { tcpEStatsStackEntry 24 }
tcpEStatsStackOtherReductions OBJECT-TYPE
   SYNTAX ZeroBasedCounter32
   MAX-ACCESS read-only
   STATUS
                  current
   DESCRIPTION
       "The number of congestion window reductions made as a result
       of anything other than AIMD congestion control algorithms.
       Examples of non-multiplicative window reductions include
       Congestion Window Validation [RFC2861] and experimental
       algorithms such as Vegas [Bra94].
```

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All window reductions MUST be counted as either tcpEStatsPerfCongSignals or tcpEStatsStackOtherReductions." REFERENCE "RFC 2861, TCP Congestion Window Validation" ::= { tcpEStatsStackEntry 25 } tcpEStatsStackCongOverCount OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only SINIAA MAX-ACCESS STATUS current DESCRIPTION "The number of congestion events that were 'backed out' of the congestion control state machine such that the congestion window was restored to a prior value. This can happen due to the Eifel algorithm [RFC3522] or other algorithms that can be used to detect and cancel spurious invocations of the Fast Retransmit Algorithm. Although it may be feasible to undo the effects of spurious invocation of the Fast Retransmit congestion events cannot easily be backed out of tcpEStatsPerfCongSignals and tcpEStatsPathPreCongSumCwnd, etc." REFERENCE "RFC 3522, The Eifel Detection Algorithm for TCP" ::= { tcpEStatsStackEntry 26 } tcpEStatsStackFastRetran OBJECT-TYPE SYNTAXZeroBasedCounter32MAX-ACCESSread-onlyCTATUCread-only current STATUS DESCRIPTION "The number of invocations of the Fast Retransmit algorithm." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsStackEntry 27 } tcpEStatsStackSubsequentTimeouts OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of times the retransmit timeout has expired after the RTO has been doubled. See Section 5.5 of RFC 2988." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsStackEntry 28 }

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tcpEStatsStackCurTimeoutCount OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current DESCRIPTION "The current number of times the retransmit timeout has expired without receiving an acknowledgment for new data. tcpEStatsStackCurTimeoutCount is reset to zero when new data is acknowledged and incremented for each invocation of Section 5.5 of RFC 2988." REFERENCE "RFC 2988, Computing TCP's Retransmission Timer" ::= { tcpEStatsStackEntry 29 } tcpEStatsStackAbruptTimeouts OBJECT-TYPE ZeroBasedCounter32 SYNTAX MAX-ACCESS read-only current STATUS DESCRIPTION "The number of timeouts that occurred without any immediately preceding duplicate acknowledgments or other indications of congestion. Abrupt Timeouts indicate that the path lost an entire window of data or acknowledgments. Timeouts that are preceded by duplicate acknowledgments or other congestion signals (e.g., ECN) are not counted as abrupt, and might have been avoided by a more sophisticated Fast Retransmit algorithm." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsStackEntry 30 } tcpEStatsStackSACKsRcvd OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of SACK options received." REFERENCE "RFC 2018, TCP Selective Acknowledgement Options" ::= { tcpEStatsStackEntry 31 } tcpEStatsStackSACKBlocksRcvd OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of SACK blocks received (within SACK options)."

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REFERENCE "RFC 2018, TCP Selective Acknowledgement Options" ::= { tcpEStatsStackEntry 32 } tcpEStatsStackSendStall OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only STATUS current DESCRIPTION "The number of interface stalls or other sender local resource limitations that are treated as congestion signals." ::= { tcpEStatsStackEntry 33 } tcpEStatsStackDSACKDups OBJECT-TYPE SYNTAX ZeroBasedCounter32 MAX-ACCESS read-only current STATUS DESCRIPTION "The number of duplicate segments reported to the local host by D-SACK blocks." REFERENCE "RFC 2883, An Extension to the Selective Acknowledgement (SACK) Option for TCP" ::= { tcpEStatsStackEntry 34 } _ _ -- The following optional objects instrument path MTU -- discovery. _ _ tcpEStatsStackMaxMSS OBJECT-TYPE SYNTAX Gauge32 UNITS "octets "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The maximum MSS, in octets." REFERENCE "RFC 1191, Path MTU discovery" ::= { tcpEStatsStackEntry 35 } tcpEStatsStackMinMSS OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION

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"The minimum MSS, in octets." REFERENCE "RFC 1191, Path MTU discovery" ::= { tcpEStatsStackEntry 36 } _ _ -- The following optional initial value objects are useful for -- conformance testing instruments on application progress and -- consumed network resources. _ _ tcpEStatsStackSndInitial OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS STATUS read-only current DESCRIPTION "Initial send sequence number. Note that by definition tcpEStatsStackSndInitial never changes for a given connection." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsStackEntry 37 } tcpEStatsStackRecInitial OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS read-only STATUS current DESCRIPTION "Initial receive sequence number. Note that by definition tcpEStatsStackRecInitial never changes for a given connection." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsStackEntry 38 } _ _ -- The following optional objects instrument the senders -- buffer usage, including any buffering in the application -- interface to TCP and the retransmit queue. All 'buffer -- memory' instruments are assumed to include OS data -- structure overhead. _ _ tcpEStatsStackCurRetxQueue OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current STATUS current Mathis, et al. Standards Track [Page 52]

DESCRIPTION "The current number of octets of data occupying the retransmit queue." ::= { tcpEStatsStackEntry 39 } tcpEStatsStackMaxRetxQueue OBJECT-TYPE SYNTAX Gauge32 "octets" UNITS MAX-ACCESS read-only STATUS current STATUS current DESCRIPTION "The maximum number of octets of data occupying the retransmit queue." ::= { tcpEStatsStackEntry 40 } tcpEStatsStackCurReasmQueue OBJECT-TYPE SYNTAX Gauge32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The current number of octets of sequence space spanned by the reassembly queue. This is generally the difference between rcv.nxt and the sequence number of the right most edge of the reassembly queue." ::= { tcpEStatsStackEntry 41 } tcpEStatsStackMaxReasmQueue OBJECT-TYPE SYNTAX Gauge32 MAX-ACCESS read-only STATUS current current STATUS DESCRIPTION "The maximum value of tcpEStatsStackCurReasmQueue" ::= { tcpEStatsStackEntry 42 } ___ -- Statistics for diagnosing interactions between -- applications and TCP. _ _ tcpEStatsAppTable OBJECT-TYPE SYNTAX SEQUENCE OF TCPEStatsAppEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "This table contains objects that are useful for determining if the application using TCP is

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```
limiting TCP performance.
         Entries are retained in this table for the number of
         seconds indicated by the tcpEStatsConnTableLatency
         object, after the TCP connection first enters the closed
         state."
    ::= { tcpEStats 6 }
tcpEStatsAppEntry OBJECT-TYPE
    SYNTAX TCpEStatsAppEntry
    MAX-ACCESS not-accessible
    STATUS current
    DESCRIPTION
         "Each entry in this table has information about the
         characteristics of each active and recently closed TCP
         connection."
   INDEX { tcpEStatsConnectIndex }
   ::= { tcpEStatsAppTable 1 }
TcpEStatsAppEntry ::= SEQUENCE {
         tcpEStatsAppSndUna
                                                  Counter32,
         tcpEStatsAppSndNxt
                                                  Unsigned32,
         Counter32,tcpEStatsAppThruOctetsAckedZeroBasedCounter32,tcpEStatsAppHCThruOctetsAckedZeroBasedCounter64,tcpEStatsAppRcvNxtCounter22tcpEStatsAppRcvNxtCounter22
         tcpEStatsAppRcvNxtcounters2,tcpEStatsAppThruOctetsReceivedZeroBasedCounter32,tcpEStatsAppHCThruOctetsReceivedZeroBasedCounter64,tcpEStatsAppCurAppWQueueGauge32,
                                        Gauge32,
Gauge32,
Cauge32
         tcpEStatsAppMaxAppWQueue
tcpEStatsAppCurAppRQueue
                                                 Gauge32,
         tcpEStatsAppMaxAppRQueue
                                                 Gauge32
    }
_ _
-- The following objects provide throughput statistics for the
-- connection including sequence numbers and elapsed
-- application data. These permit direct observation of the
-- applications progress, in terms of elapsed data delivery
-- and elapsed time.
_ _
tcpEStatsAppSndUna OBJECT-TYPE
    SYNTAX Counter32
    MAX-ACCESS read-only
    STATUS
                     current
    DESCRIPTION
```

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"The value of SND.UNA, the oldest unacknowledged sequence number. Note that SND.UNA is a TCP state variable that is congruent to Counter32 semantics." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsAppEntry 1 } tcpEStatsAppSndNxt OBJECT-TYPE SYNTAX Unsigned32 MAX-ACCESS read-only MAX-ACCESS STATUS current DESCRIPTION "The value of SND.NXT, the next sequence number to be sent. Note that tcpEStatsAppSndNxt is not monotonic (and thus not a counter) because TCP sometimes retransmits lost data by pulling tcpEStatsAppSndNxt back to the missing data." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsAppEntry 2 } tcpEStatsAppSndMax OBJECT-TYPE Counter32 read-only SYNTAX MAX-ACCESS STATUS current DESCRIPTION "The farthest forward (right most or largest) SND.NXT value. Note that this will be equal to tcpEStatsAppSndNxt except when tcpEStatsAppSndNxt is pulled back during recovery." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsAppEntry 3 } tcpEStatsAppThruOctetsAcked OBJECT-TYPE SYNTAX ZeroBasedCounter32 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets for which cumulative acknowledgments have been received. Note that this will be the sum of changes to tcpEStatsAppSndUna." ::= { tcpEStatsAppEntry 4 } tcpEStatsAppHCThruOctetsAcked OBJECT-TYPE SYNTAX ZeroBasedCounter64 UNITS "octets"

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MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets for which cumulative acknowledgments have been received, on systems that can receive more than 10 million bits per second. Note that this will be the sum of changes in tcpEStatsAppSndUna." ::= { tcpEStatsAppEntry 5 } tcpEStatsAppRcvNxt OBJECT-TYPE Counter32 S read-only SYNTAX MAX-ACCESS STATUS current DESCRIPTION "The value of RCV.NXT. The next sequence number expected on an incoming segment, and the left or lower edge of the receive window. Note that RCV.NXT is a TCP state variable that is congruent to Counter32 semantics." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsAppEntry 6 } tcpEStatsAppThruOctetsReceived OBJECT-TYPE SYNTAX ZeroBasedCounter32 "octets" UNITS MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets for which cumulative acknowledgments have been sent. Note that this will be the sum of changes to tcpEStatsAppRcvNxt." ::= { tcpEStatsAppEntry 7 } tcpEStatsAppHCThruOctetsReceived OBJECT-TYPE SYNTAX ZeroBasedCounter64 UNITS "octets" MAX-ACCESS read-only STATUS current DESCRIPTION "The number of octets for which cumulative acknowledgments have been sent, on systems that can transmit more than 10 million bits per second. Note that this will be the sum of changes in tcpEStatsAppRcvNxt." ::= { tcpEStatsAppEntry 8 } tcpEStatsAppCurAppWQueue OBJECT-TYPE

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SYNTAX	Gauge32
UNITS	"octets"
MAX-ACCESS	read-only
STATUS	current
DESCRIPTION	
by TCP, per SND.NXT or (and SND.NX available o (rwin). Th transmissio	
tcpEStatsAppMaxAppW	
SYNTAX	Gauge32
UNITS	"octets"
MAX-ACCESS	read-only
STATUS	current
DESCRIPTION	n number of octets of application data buffered
value of to be used to state (sugg (suggesting	nding first transmission. This is the maximum cpEStatsAppCurAppWQueue. This pair of objects can determine if insufficient queued data is steady gesting insufficient queue space) or transient g insufficient application performance or CPU load or scheduler latency)." sAppEntry 12 }
tcpEStatsAppCurAppH	Queue OBJECT-TYPE
SYNTAX	Gauge32
UNITS	"octets"
MAX-ACCESS	read-only
STATUS	current
	when it is l
tcpEStatsAppMaxAppH	Queue OBJECT-TYPE
	Gauge32
SYNTAX	
SYNTAX UNITS	"octets"
SYNTAX UNITS MAX-ACCESS	read-only
SYNTAX UNITS	

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"The maximum number of octets of application data that has been acknowledged by TCP but not yet delivered to the application." ::= { tcpEStatsAppEntry 14 } -- Controls for Tuning TCP _ _ tcpEStatsTuneTable OBJECT-TYPE SYNTAX SEQUENCE OF TCPEStatsTuneEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "This table contains per-connection controls that can be used to work around a number of common problems that plague TCP over some paths. All can be characterized as limiting the growth of the congestion window so as to prevent TCP from overwhelming some component in the path. Entries are retained in this table for the number of seconds indicated by the tcpEStatsConnTableLatency object, after the TCP connection first enters the closed state." ::= { tcpEStats 7 } tcpEStatsTuneEntry OBJECT-TYPE SYNTAX TcpEStatsTuneEntry MAX-ACCESS not-accessible STATUS current DESCRIPTION "Each entry in this table is a control that can be used to place limits on each active TCP connection." INDEX { tcpEStatsConnectIndex } ::= { tcpEStatsTuneTable 1 } TcpEStatsTuneEntry ::= SEQUENCE { Unsigned32, Unsigned32, Unsigned32, tcpEStatsTuneLimCwnd tcpEStatsTuneLimCwnd tcpEStatsTuneLimSsthresh tcpEStatsTuneLimRwin tcpEStatsTuneLimRwin tcpEStatsTuneLimMSS Unsigned32 } tcpEStatsTuneLimCwnd OBJECT-TYPE Unsigned32 SYNTAX

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UNITS "octets" UNIIS MAX-ACCESS read-write STATUS current DESCRIPTION "A control to set the maximum congestion window that may be used, in octets." REFERENCE "RFC 2581, TCP Congestion Control" ::= { tcpEStatsTuneEntry 1 } tcpEStatsTuneLimSsthresh OBJECT-TYPE SYNTAX Unsigned32 UNITS "octets" MAX-ACCESS read-write STATUS current DESCRIPTION "A control to limit the maximum queue space (in octets) that this TCP connection is likely to occupy during slowstart. It can be implemented with the algorithm described in RFC 3742 by setting the max_ssthresh parameter to twice tcpEStatsTuneLimSsthresh. This algorithm can be used to overcome some TCP performance problems over network paths that do not have sufficient buffering to withstand the bursts normally present during slowstart." REFERENCE "RFC 3742, Limited Slow-Start for TCP with Large Congestion Windows" ::= { tcpEStatsTuneEntry 2 } tcpEStatsTuneLimRwin OBJECT-TYPE SYNTAX Unsigned32 UNITS "octets" MAX-ACCESS read-write STATUS current DESCRIPTION "A control to set the maximum window advertisement that may be sent, in octets." REFERENCE "RFC 793, Transmission Control Protocol" ::= { tcpEStatsTuneEntry 3 } tcpEStatsTuneLimMSS OBJECT-TYPE SYNTAX Unsigned32 UNITS "octets" MAX-ACCESS read-write

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```
STATUS
                current
   DESCRIPTION
     "A control to limit the maximum segment size in octets, that
      this TCP connection can use."
   REFERENCE
     "RFC 1191, Path MTU discovery"
   ::= { tcpEStatsTuneEntry 4 }
-- TCP Extended Statistics Notifications Group
_ _
tcpEStatsEstablishNotification NOTIFICATION-TYPE
   OBJECTS {
               tcpEStatsConnectIndex
             }
         current
   STATUS
   DESCRIPTION
      "The indicated connection has been accepted
      (or alternatively entered the established state)."
   ::= { tcpEStatsNotifications 1 }
tcpEStatsCloseNotification NOTIFICATION-TYPE
   OBJECTS {
               tcpEStatsConnectIndex
             }
   STATUS current
   DESCRIPTION
      "The indicated connection has left the
      established state"
   ::= { tcpEStatsNotifications 2 }
_ _
-- Conformance Definitions
_ _
  tcpEStatsCompliances OBJECT IDENTIFIER
      ::= { tcpEStatsConformance 1 }
  tcpEStatsGroups OBJECT IDENTIFIER
      ::= { tcpEStatsConformance 2 }
-- Compliance Statements
 tcpEStatsCompliance MODULE-COMPLIANCE
```

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```
STATUS current
DESCRIPTION
    "Compliance statement for all systems that implement TCP
    extended statistics."
MODULE -- this module
    MANDATORY-GROUPS {
                       tcpEStatsListenerGroup,
                       tcpEStatsConnectIdGroup,
                       tcpEStatsPerfGroup,
                       tcpEStatsPathGroup,
                       tcpEStatsStackGroup,
                       tcpEStatsAppGroup
                     }
    GROUP tcpEStatsListenerHCGroup
    DESCRIPTION
        "This group is mandatory for all systems that can
         wrap the values of the 32-bit counters in
         tcpEStatsListenerGroup in less than one hour."
    GROUP tcpEStatsPerfOptionalGroup
    DESCRIPTION
        "This group is optional for all systems."
    GROUP tcpEStatsPerfHCGroup
    DESCRIPTION
        "This group is mandatory for systems that can
        wrap the values of the 32-bit counters in
        tcpEStatsPerfGroup in less than one hour.
        Note that any system that can attain 10 Mb/s
        can potentially wrap 32-Bit Octet counters in
        under one hour."
    GROUP tcpEStatsPathOptionalGroup
    DESCRIPTION
        "This group is optional for all systems."
    GROUP tcpEStatsPathHCGroup
    DESCRIPTION
        "This group is mandatory for systems that can
        wrap the values of the 32-bit counters in
        tcpEStatsPathGroup in less than one hour.
        Note that any system that can attain 10 \,\mathrm{Mb/s}
        can potentially wrap 32-Bit Octet counters in
        under one hour."
    GROUP tcpEStatsStackOptionalGroup
```

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DESCRIPTION "This group is optional for all systems." GROUP tcpEStatsAppHCGroup DESCRIPTION "This group is mandatory for systems that can wrap the values of the 32-bit counters in tcpEStatsStackGroup in less than one hour. Note that any system that can attain 10 Mb/s can potentially wrap 32-Bit Octet counters in under one hour." GROUP tcpEStatsAppOptionalGroup DESCRIPTION "This group is optional for all systems." GROUP tcpEStatsTuneOptionalGroup DESCRIPTION "This group is optional for all systems." GROUP tcpEStatsNotificationsGroup DESCRIPTION "This group is optional for all systems." GROUP tcpEStatsNotificationsCtlGroup DESCRIPTION "This group is mandatory for systems that include the tcpEStatsNotificationGroup." ::= { tcpEStatsCompliances 1 } -- Units of Conformance tcpEStatsListenerGroup OBJECT-GROUP OBJECTS { tcpEStatsListenerTableLastChange, tcpEStatsListenerStartTime, tcpEStatsListenerSynRcvd, tcpEStatsListenerInitial, tcpEStatsListenerEstablished, tcpEStatsListenerAccepted, tcpEStatsListenerExceedBacklog, tcpEStatsListenerCurConns, tcpEStatsListenerMaxBacklog, tcpEStatsListenerCurBacklog,

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_ _

_ _

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tcpEStatsListenerCurEstabBacklog } STATUS current DESCRIPTION "The tcpEStatsListener group includes objects that provide valuable statistics and debugging information for TCP Listeners." ::= { tcpEStatsGroups 1 } tcpEStatsListenerHCGroup OBJECT-GROUP OBJECTS { tcpEStatsListenerHCSynRcvd, tcpEStatsListenerHCInitial, tcpEStatsListenerHCEstablished, tcpEStatsListenerHCAccepted, tcpEStatsListenerHCExceedBacklog } STATUS current DESCRIPTION "The tcpEStatsListenerHC group includes 64-bit counters in tcpEStatsListenerTable." ::= { tcpEStatsGroups 2 } tcpEStatsConnectIdGroup OBJECT-GROUP OBJECTS { tcpEStatsConnTableLatency, tcpEStatsConnectIndex } STATUS current DESCRIPTION "The tcpEStatsConnectId group includes objects that identify TCP connections and control how long TCP connection entries are retained in the tables." ::= { tcpEStatsGroups 3 } tcpEStatsPerfGroup OBJECT-GROUP OBJECTS { tcpEStatsPerfSegsOut, tcpEStatsPerfDataSegsOut, tcpEStatsPerfDataOctetsOut, tcpEStatsPerfSegsRetrans, tcpEStatsPerfOctetsRetrans, tcpEStatsPerfSegsIn, tcpEStatsPerfDataSegsIn, tcpEStatsPerfDataOctetsIn, tcpEStatsPerfElapsedSecs, tcpEStatsPerfElapsedMicroSecs, tcpEStatsPerfStartTimeStamp, tcpEStatsPerfCurMSS, tcpEStatsPerfPipeSize, tcpEStatsPerfMaxPipeSize, tcpEStatsPerfSmoothedRTT, tcpEStatsPerfCurRTO,

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```
tcpEStatsPerfCongSignals, tcpEStatsPerfCurCwnd,
          tcpEStatsPerfCurSsthresh, tcpEStatsPerfTimeouts,
          tcpEStatsPerfCurRwinSent,
          tcpEStatsPerfMaxRwinSent,
          tcpEStatsPerfZeroRwinSent,
          tcpEStatsPerfCurRwinRcvd,
          tcpEStatsPerfMaxRwinRcvd,
          tcpEStatsPerfZeroRwinRcvd
     }
     STATUS current
    DESCRIPTION
          "The tcpEStatsPerf group includes those objects that
          provide basic performance data for a TCP connection."
  ::= { tcpEStatsGroups 4 }
tcpEStatsPerfOptionalGroup OBJECT-GROUP
    OBJECTS {
          tcpEStatsPerfSndLimTransRwin,
          tcpEStatsPerfSndLimTransCwnd,
          tcpEStatsPerfSndLimTransSnd,
          tcpEStatsPerfSndLimTimeRwin,
          tcpEStatsPerfSndLimTimeCwnd,
          tcpEStatsPerfSndLimTimeSnd
     STATUS current
    DESCRIPTION
          "The tcpEStatsPerf group includes those objects that
          provide basic performance data for a TCP connection."
  ::= { tcpEStatsGroups 5 }
tcpEStatsPerfHCGroup OBJECT-GROUP
    OBJECTS {
          tcpEStatsPerfHCDataOctetsOut,
          tcpEStatsPerfHCDataOctetsIn
     }
     STATUS current
    DESCRIPTION
          "The tcpEStatsPerfHC group includes 64-bit
          counters in the tcpEStatsPerfTable."
  ::= { tcpEStatsGroups 6 }
tcpEStatsPathGroup OBJECT-GROUP
    OBJECTS {
          tcpEStatsControlPath,
          tcpEStatsPathRetranThresh,
          tcpEStatsPathNonRecovDAEpisodes,
          tcpEStatsPathSumOctetsReordered,
```

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tcpEStatsPathNonRecovDA } STATUS current DESCRIPTION "The tcpEStatsPath group includes objects that control the creation of the tcpEStatsPathTable, and provide information about the path for each TCP connection." ::= { tcpEStatsGroups 7 } tcpEStatsPathOptionalGroup OBJECT-GROUP OBJECTS { tcpEStatsPathSampleRTT, tcpEStatsPathRTTVar, tcpEStatsPathMaxRTT, tcpEStatsPathMinRTT, tcpEStatsPathSumRTT, tcpEStatsPathCountRTT, tcpEStatsPathMaxRTO, tcpEStatsPathMinRTO, tcpEStatsPathIpTtl, tcpEStatsPathIpTosIn, tcpEStatsPathIpTosOut, tcpEStatsPathPreCongSumCwnd, tcpEStatsPathPreCongSumRTT, tcpEStatsPathPostCongSumRTT, tcpEStatsPathPostCongCountRTT, tcpEStatsPathECNsignals, tcpEStatsPathDupAckEpisodes, tcpEStatsPathRcvRTT, tcpEStatsPathDupAcksOut, tcpEStatsPathCERcvd, tcpEStatsPathECESent STATUS current DESCRIPTION "The tcpEStatsPath group includes objects that provide additional information about the path for each TCP connection." ::= { tcpEStatsGroups 8 } tcpEStatsPathHCGroup OBJECT-GROUP OBJECTS { tcpEStatsPathHCSumRTT STATUS current DESCRIPTION "The tcpEStatsPathHC group includes 64-bit counters in the tcpEStatsPathTable." ::= { tcpEStatsGroups 9 } tcpEStatsStackGroup OBJECT-GROUP OBJECTS { tcpEStatsControlStack, tcpEStatsStackActiveOpen, tcpEStatsStackMSSSent,

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tcpEStatsStackMSSRcvd, tcpEStatsStackWinScaleSent, tcpEStatsStackWinScaleRcvd, tcpEStatsStackTimeStamps, tcpEStatsStackECN, tcpEStatsStackWillSendSACK, tcpEStatsStackWillUseSACK, tcpEStatsStackState, tcpEStatsStackNagle, tcpEStatsStackMaxSsCwnd, tcpEStatsStackMaxCaCwnd, tcpEStatsStackMaxSsthresh, tcpEStatsStackMinSsthresh, tcpEStatsStackInRecovery, tcpEStatsStackDupAcksIn, tcpEStatsStackSpuriousFrDetected, tcpEStatsStackSpuriousRtoDetected } STATUS current DESCRIPTION "The tcpEStatsConnState group includes objects that control the creation of the tcpEStatsStackTable, and provide information about the operation of algorithms used within TCP." ::= { tcpEStatsGroups 10 } tcpEStatsStackOptionalGroup OBJECT-GROUP OBJECTS { tcpEStatsStackSoftErrors, tcpEStatsStackSoftErrorReason, tcpEStatsStackSlowStart, tcpEStatsStackCongAvoid, tcpEStatsStackOtherReductions, tcpEStatsStackCongOverCount, tcpEStatsStackFastRetran, tcpEStatsStackSubsequentTimeouts, tcpEStatsStackCurTimeoutCount, tcpEStatsStackAbruptTimeouts, tcpEStatsStackSACKsRcvd, tcpEStatsStackSACKBlocksRcvd, tcpEStatsStackSendStall, tcpEStatsStackDSACKDups, tcpEStatsStackMaxMSS, tcpEStatsStackMinMSS, tcpEStatsStackSndInitial, tcpEStatsStackRecInitial, tcpEStatsStackCurRetxQueue, tcpEStatsStackMaxRetxQueue, tcpEStatsStackCurReasmQueue, tcpEStatsStackMaxReasmQueue STATUS current DESCRIPTION "The tcpEStatsConnState group includes objects that provide additional information about the operation of algorithms used within TCP."

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::= { tcpEStatsGroups 11 }

```
tcpEStatsAppGroup OBJECT-GROUP
       OBJECTS {
           tcpEStatsControlApp,
            tcpEStatsAppSndUna, tcpEStatsAppSndNxt,
            tcpEStatsAppSndMax, tcpEStatsAppThruOctetsAcked,
            tcpEStatsAppRcvNxt,
            tcpEStatsAppThruOctetsReceived
       }
       STATUS current
      DESCRIPTION
            "The tcpEStatsConnState group includes objects that
           control the creation of the tcpEStatsAppTable,
           and provide information about the operation of
            algorithms used within TCP."
    ::= { tcpEStatsGroups 12 }
tcpEStatsAppHCGroup OBJECT-GROUP
      OBJECTS {
            tcpEStatsAppHCThruOctetsAcked,
            tcpEStatsAppHCThruOctetsReceived
       }
       STATUS current
      DESCRIPTION
            "The tcpEStatsStackHC group includes 64-bit
            counters in the tcpEStatsStackTable."
    ::= { tcpEStatsGroups 13 }
 tcpEStatsAppOptionalGroup OBJECT-GROUP
      OBJECTS {
           tcpEStatsAppCurAppWQueue,
           tcpEStatsAppMaxAppWQueue,
           tcpEStatsAppCurAppRQueue,
           tcpEStatsAppMaxAppRQueue
       STATUS current
      DESCRIPTION
            "The tcpEStatsConnState group includes objects that
           provide additional information about how applications
            are interacting with each TCP connection."
    ::= { tcpEStatsGroups 14 }
 tcpEStatsTuneOptionalGroup OBJECT-GROUP
      OBJECTS {
           tcpEStatsControlTune,
            tcpEStatsTuneLimCwnd, tcpEStatsTuneLimSsthresh,
            tcpEStatsTuneLimRwin, tcpEStatsTuneLimMSS
```

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```
}
     STATUS current
    DESCRIPTION
         "The tcpEStatsConnState group includes objects that
         control the creation of the tcpEStatsConnectionTable,
         which can be used to set tuning parameters
         for each TCP connection."
  ::= { tcpEStatsGroups 15 }
tcpEStatsNotificationsGroup
                               NOTIFICATION-GROUP
    NOTIFICATIONS {
                  tcpEStatsEstablishNotification,
                  tcpEStatsCloseNotification
     }
    STATUS current
    DESCRIPTION
         "Notifications sent by a TCP extended statistics agent."
  ::= { tcpEStatsGroups 16 }
tcpEStatsNotificationsCtlGroup OBJECT-GROUP
    OBJECTS {
                  tcpEStatsControlNotify
     }
    STATUS current
    DESCRIPTION
         "The tcpEStatsNotificationsCtl group includes the
          object that controls the creation of the events
         in the tcpEStatsNotificationsGroup."
  ::= { tcpEStatsGroups 17 }
```

END

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5. Security Considerations

There are a number of management objects defined in this MIB module with a MAX-ACCESS clause of read-write and/or read-create. Such objects may be considered sensitive or vulnerable in some network environments. The support for SET operations in a non-secure environment without proper protection can have a negative effect on network operations. These are the tables and objects and their sensitivity/vulnerability:

- * Changing tcpEStatsConnTableLatency or any of the control objects in the tcpEStatsControl group (tcpEStatsControlPath, tcpEStatsControlStack, tcpEStatsControlApp, tcpEStatsControlTune) may affect the correctness of other management applications accessing this MIB. Generally, local policy should only permit limited write access to these controls (e.g., only by one management station or only during system configuration).
- * The objects in the tcpEStatsControlTune group (tcpEStatsTuneLimCwnd, tcpEStatsTuneLimSsthresh, tcpEStatsTuneLimRwin) can be used to limit resources consumed by TCP connections or to limit TCP throughput. An attacker might manipulate these objects to reduce performance to levels below the minimum acceptable for a particular application.

Some of the readable objects in this MIB module (i.e., objects with a MAX-ACCESS other than not-accessible) may be considered sensitive or vulnerable in some network environments. It is thus important to control even GET and/or NOTIFY access to these objects and possibly to even encrypt the values of these objects when sending them over the network via SNMP. These are the tables and objects and their sensitivity/vulnerability:

- * All objects which expose TCP sequence numbers (tcpEStatsAppSndUna, tcpEStatsAppSndNxt, tcpEStatsAppSndMax, tcpEStatsStackSndInitial, tcpEStatsAppRcvNxt, and tcpEStatsStackRecInitial) might make it easier for an attacker to forge in sequence TCP segments to disrupt TCP connections.
- * Nearly all objects in this (or any other) MIB may be used to estimate traffic volumes, which may reveal unanticipated information about an organization to the outside world.

SNMP versions prior to SNMPv3 did not include adequate security. Even if the network itself is secure (for example by using IPsec), even then, there is no control as to who on the secure network is allowed to access and GET/SET (read/change/create/delete) the objects in this MIB module.

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It is RECOMMENDED that implementers consider the security features as provided by the SNMPv3 framework (see [RFC3410], section 8), including full support for the SNMPv3 cryptographic mechanisms (for authentication and privacy).

Further, deployment of SNMP versions prior to SNMPv3 is NOT RECOMMENDED. Instead, it is RECOMMENDED to deploy SNMPv3 and to enable cryptographic security. It is then a customer/operator responsibility to ensure that the SNMP entity giving access to an instance of this MIB module is properly configured to give access to the objects only to those principals (users) that have legitimate rights to indeed GET or SET (change/create/delete) them.

6. IANA Considerations

The MIB module in this document uses the following IANA-assigned OBJECT IDENTIFIER values recorded in the SMI Numbers registry:

Descriptor	OBJECT IDENTIFIER value
tcpEStatsMIB	{ mib-2 156 }

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9. Contributors

The following people contributed text that was incorporated into this document:

Jon Saperia <saperia@jdscons.com> converted Web100 internal documentation into a true MIB.

Some of the objects in this document were moved from an early version of the TCP-MIB by Bill Fenner, et al.

Some of the object descriptions are based on an earlier unpublished document by Jeff Semke.

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